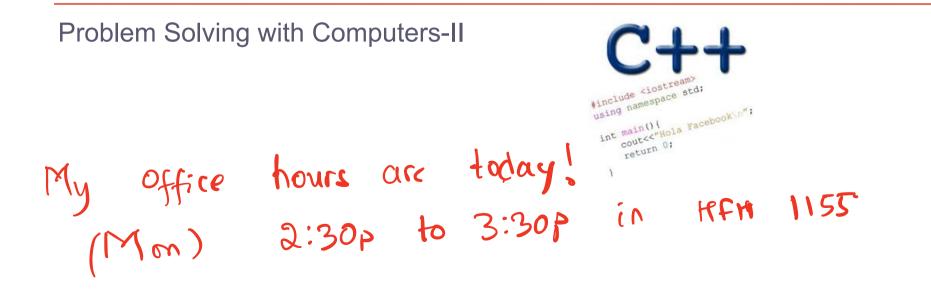
BINARY SEARCH TREES

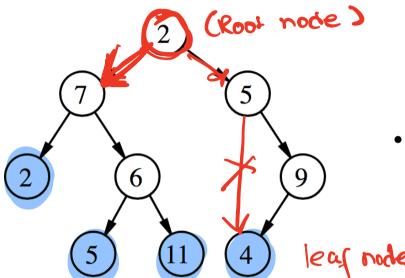


Binary Search

- Binary search. Given value and sorted array a[], find index i such that a[i] = value, or report that no such index exists.
- Invariant. Algorithm maintains a[lo] ≤ value ≤ a[hi].
- Ex. Binary search for 33.

6	13	14	25	33	43	51	53	64	72	84	93	95	96	97
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14
Î														Î
lo														hi



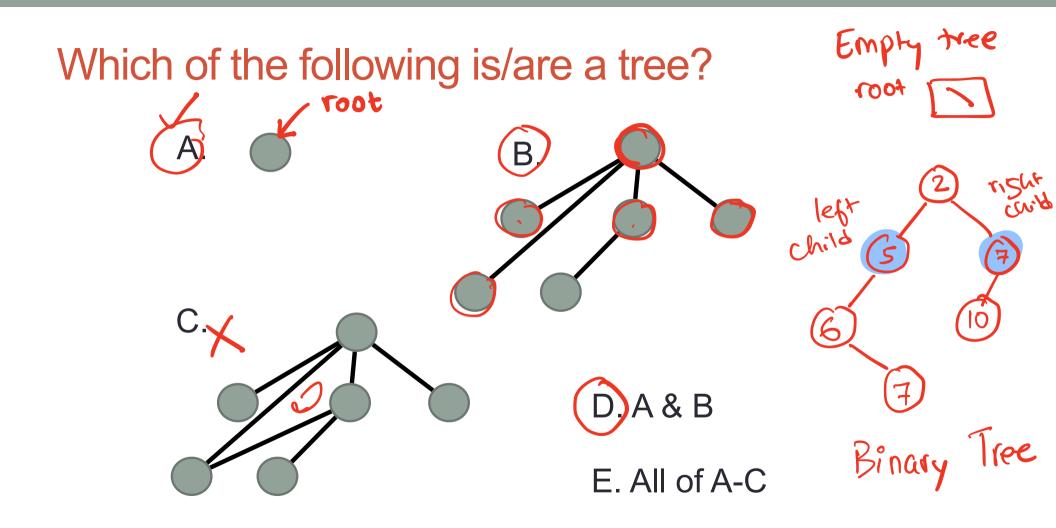


A tree has following general properties:

- One node is distinguished as a **root**;
- Every node (exclude a root) is connected by a directed edge *from* exactly one other node;
 - A direction is: *parent -> children*
- Leaf node: Node that has no children

nod-e 1007 the children are 25 leap nodes 2's





Binary Search Trees

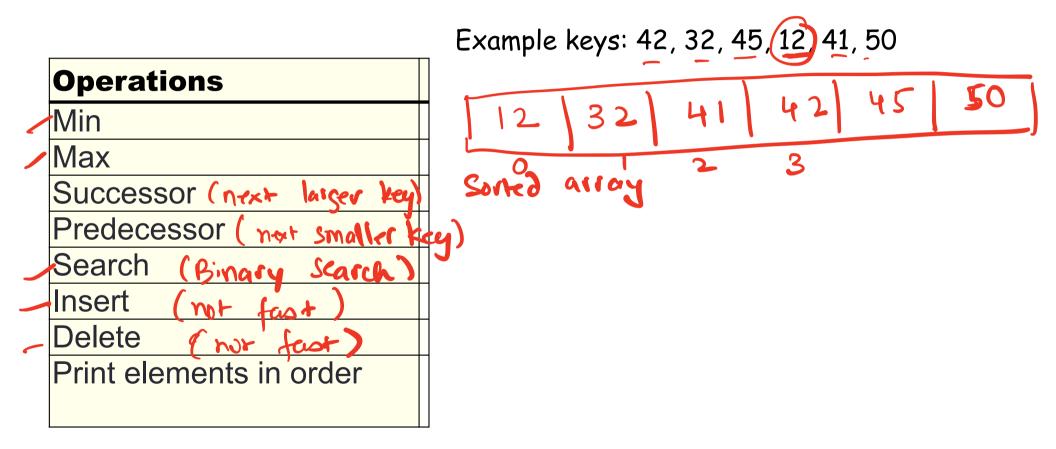
What are the operations supported?
 Scarch :

• What are the running times of these operations?

next lectures

• How do you implement the BST i.e. operations supported by it?

Operations supported by Sorted arrays and Binary Search Trees (BST)



Left run Binary Search Tree – What is it? 1e(42) 32 45 12 50 41 $T_{L}(k) \langle k \langle T_{R}(k)$

Each node:

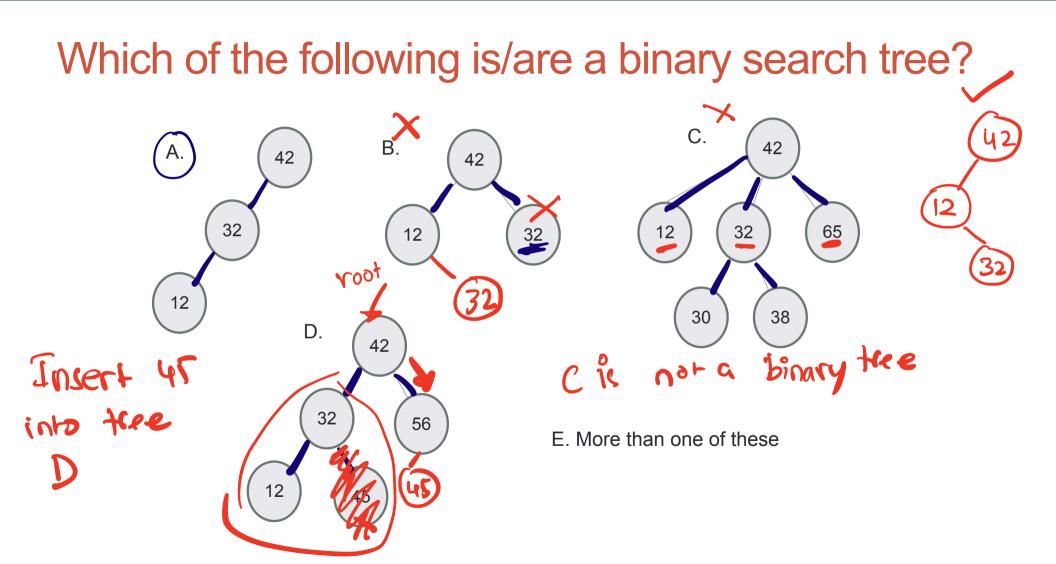
- stores a key (k)
 - has a pointer to left child, right child and parent (optional)

7

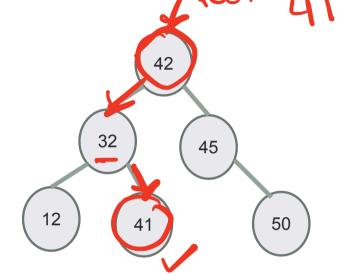
Satisfies the Search Tree Property

For any node,

Keys in node's left subtree < Node's key Node's key < Keys in node's right subtree



BSTs allow efficient search!

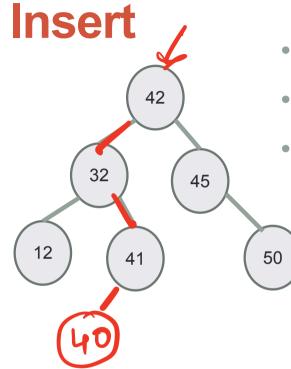


- Start at the root;
- Trace down a path by comparing **k** with the key of the current node x:
 - If the keys are equal: we have found the key
 - If **k** < key[x] search in the left subtree of x
 - If **k** > key[x] search in the right subtree of x

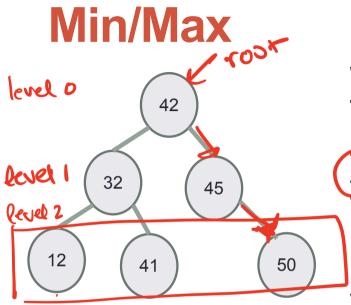
2 32 41 42 45 50



Search for 41, then search for 53



- Insert 40
- Search for the key
- Insert at the spot you expected to find it



Which of the following described the algorithm to find the maximum value in the BST?

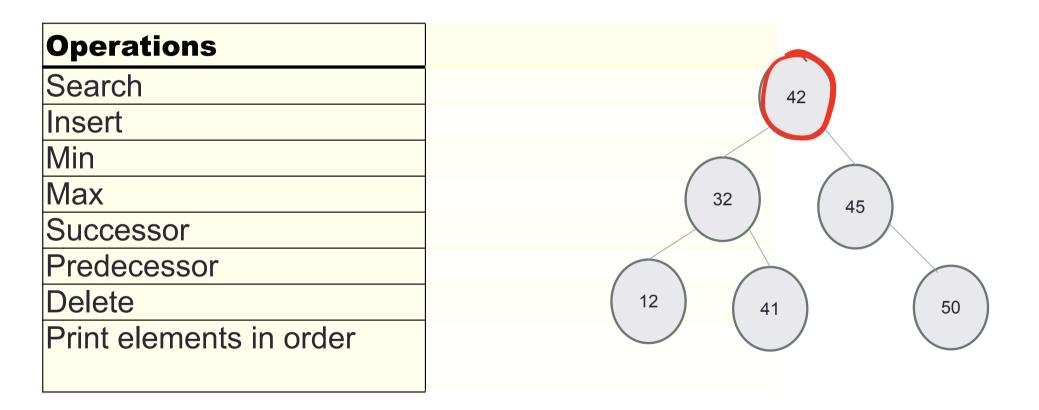
A. Follow right child pointers from the root, until a node with no right child is encountered, return that node's key

(min)

B. Follow left child pointers from the root, until a node with no left child is encountered, return that node's key

C. Traverse to the last level in the tree and traverse the tree left to right, return the key of the last node in the last level. Description is not precise enough to be an algorithm Need to define the notion of "level" & how to traverse a tevel from test to right. to traverse a tevel from test to right. Finally. even if we could carry out these steps. algo doesn't work for any BST. Look for a counterexample.

Define the BST ADT

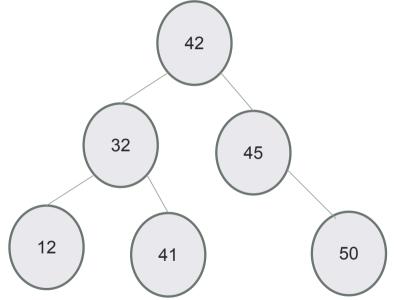


```
class BSTNode {
                                               data navi
public:
                                           RCT
                                                            parent
  BSTNode* left;
  BSTNode* right;
  BSTNode* parent;
                                                       data
  int const data;
  BSTNode ( const int e d ) : data(d) {
     left = right = parent = \partial \Re_{u} \otimes \mathcal{O}_{v};
```

13

Traversing down the tree

- Suppose n is a pointer to the root. What is the output of the following code:
 - n = n->left;
 - n = n->right;
 - cout<<n->data<<endl;</pre>
 - A. 42
 - B. 32
 - C. 12
 - D. 41
 - E. Segfault



Traversing up the tree

- Suppose n is a pointer to the node with value 50.
- What is the output of the following code:

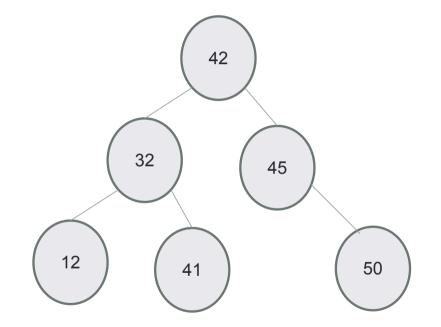
```
n = n->parent;
```

- n = n->parent;
- n = n->left;

```
cout<<n->data<<endl;</pre>
```

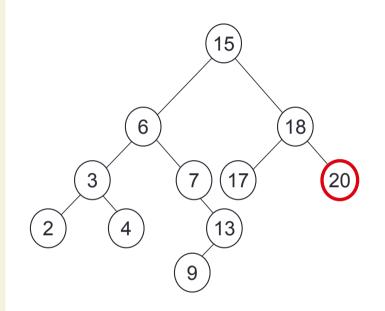
- A. 42
- B. 32
- C. 12
- D. 45

E. Segfault



Max: find the maximum key value in a BST

Alg: int BST::max()



Maximum = 20

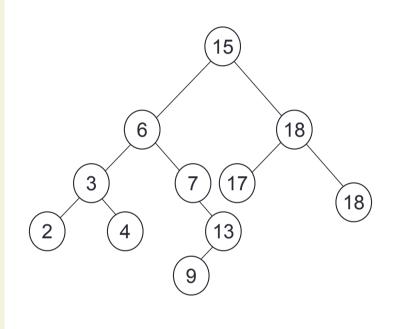
16

Min: find the minimum key value in a BST

```
Alg: int BST::min() {
```

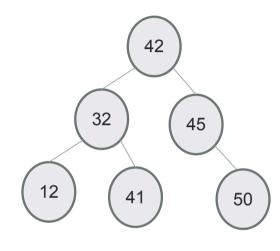
```
Start at the root.
Follow child
pointers from the root, until
a node with no left child is
encountered.
```

Return the key of that node



Min = ?

In order traversal: print elements in sorted order



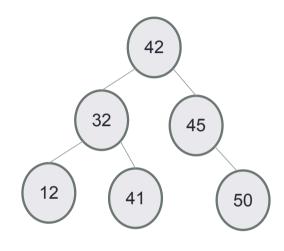
Algorithm Inorder(tree)

1. Traverse the left subtree, i.e., call Inorder(left-subtree)

2. Visit the root.

3. Traverse the right subtree, i.e., call Inorder(right-subtree)

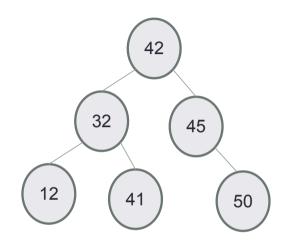
Pre-order traversal: nice way to linearize your tree!



Algorithm Preorder(tree)

- 1. Visit the root.
- 2. Traverse the left subtree, i.e., call Preorder(left-subtree)
- 3. Traverse the right subtree, i.e., call Preorder(right-subtree)

Post-order traversal: use in recursive destructors!

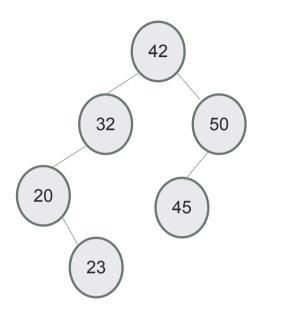


Algorithm Postorder(tree)

- 1. Traverse the left subtree, i.e., call Postorder(left-subtree)
- 2. Traverse the right subtree, i.e., call Postorder(right-subtree)

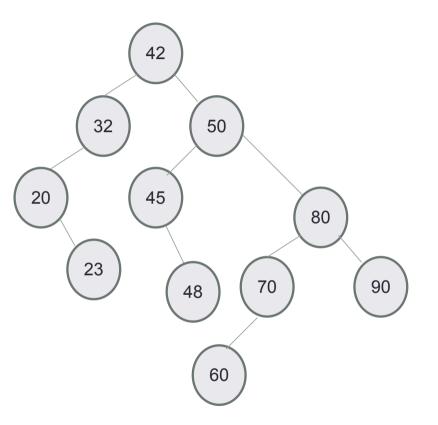
3. Visit the root.

Predecessor: Next smallest element



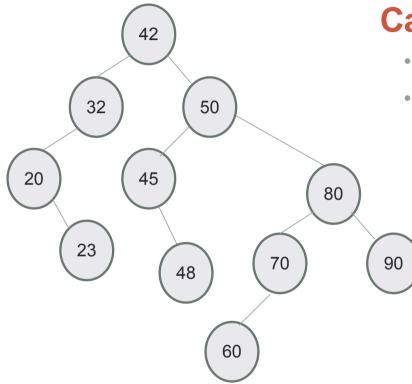
- What is the predecessor of 32?
- What is the predecessor of 45?

Successor: Next largest element



- What is the successor of 45?
- What is the successor of 50?
- What is the successor of 60?

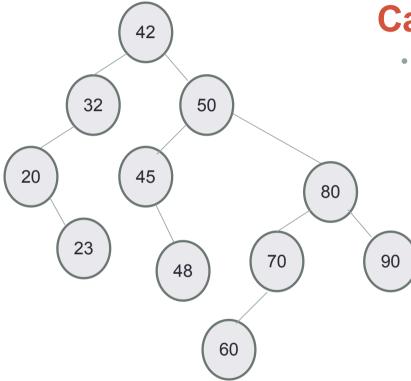
Delete: Case 1



Case 1: Node is a leaf node

- Set parent's (left/right) child pointer to null
- Delete the node

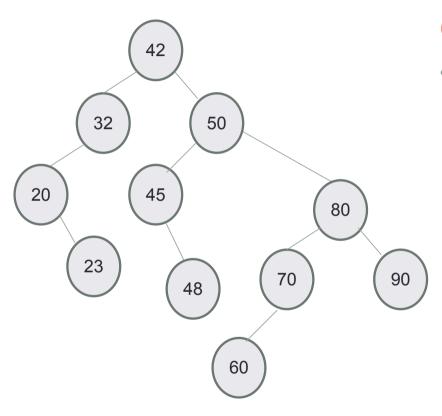
Delete: Case 2



Case 2 Node has only one child

• Replace the node by its only child

Delete: Case 3



Case 3 Node has two children

• Can we still replace the node by one of its children? Why or Why not?