# HEAPS

Problem Solving with Computers-II





## How is PA02 going?

- A. Done
- B. On track to finish
- C. Having some difficulties
- D. Just started
- E. Haven't started

#### Heaps

- Clarification
  - *heap*, the data structure is not related to *heap*, the region of memory
- What are the operations supported?
- What are the running times?

#### Heaps

Min-Heaps

Max-Heap

BST

• Insert :

- Min:
- Delete Min:
- Max
- Delete Max

#### **Applications:**

- Efficient sort
- Finding the median of a sequence of numbers
- Compression codes

Choose heap if you are doing repeated insert/delete/(min OR max) operations

#### std::priority\_queue (STL's version of heap)

A C++ priority\_queue is a generic container, and can store any data type on which an ordering can be defined: for example ints, structs (Card), pointers etc.

#include <queue>
priority\_queue<int> pq;

#### Methods:

- \* push() //insert
- \* pop() //delete max priority item
- \*top() //get max priority item
- \* empty() //returns true if the priority queue is empty
- You can extract object of highest priority in O(log N)
- To determine priority: objects in a priority queue must be comparable to each other

STL Heap implementation: Priority Queues in C++

What is the output of this code?

```
priority queue<int> pq;
pq.push(10);
pq.push(2);
pq.push(80);
cout<<pq.top();</pre>
pq.pop();
cout<<pq.top();</pre>
pq.pop();
cout<<pq.top();</pre>
pq.pop();
```

A.10 2 80 B.2 10 80 C.80 10 2 D.80 2 10 E. None of the above

#### Heaps as binary trees

- Rooted binary tree that is as complete as possible
- In a min-Heap, each node satisfies the following heap property:

key(x)<= key(children of x)</pre>





#### Heaps as binary trees

- Rooted binary tree that is as complete as possible
- In a max-Heap, each node satisfies the following heap property: key(x)>= key(children of x)



## Identifying heaps

Starting with the following min-Heap which of the following operations will result in something that is NOT a min Heap

- A. Swap the nodes 40 and 32
- B. Swap the nodes 32 and 43
- C. Swap the nodes 43 and 40
- D. Insert 50 as the left child of 45 E. C&D



#### Structure: Complete binary tree

A heap is a complete binary tree: Each level is as full as possible. Nodes on the bottom level are placed as far left as possible



### Insert 50 into a heap

- Insert key(x) in the first open slot at the last level of tree (going from left to right)
- If the heap property is not violated Done
- Else: while(key(parent(x))>key(x)) swap the key(x) with key(parent(x))



#### Insert 50, then 35, then 8



#### **Delete min**

- Replace the root with the rightmost node at the last level
- "Bubble down"- swap node with one of the children until the heap property is restored



### Under the hood of heaps

- An efficient way of implementing heaps is using vectors
- Although we think of heaps as trees, the entire tree can be efficiently represented as a vector!!

#### Implementing heaps using an array or vector



### Insert into a heap

- Insert key(x) in the first open slot at the last level of tree (going from left to right)
- If the heap property is not violated Done
- Else....

Insert the elements {12, 41, 47, 45, 32} in a min-Heap using the vector representation of the heap

Finding the "parent" of a "node" in the vector representation



For a node at index i, index of the parent is int(i-1/2)

Value	6	10	12	40	32	43	47	45	41	
Index	0	1	2	3	4	5	6	7	8	

#### Insert 50, then 35



For a node at index i, index of the parent is int(i-1/2)

Value	6	10	12	40	32	43	47	45	41	
Index	0	1	2	3	4	5	6	7	8	

#### Insert 8 into a heap

Value	6	10	12	40	32	43	47	45	41	50	35
Index	0	1	2	3	4	5	6	7	8	9	10

After inserting 8, which node is the parent of 8 ?

- A. Node 6
- B. Node 12
- C. None 43
- **D. None Node 8 will be the root**

#### **Delete min**

- Replace the root with the rightmost node at the last level
- "Bubble down"- swap node with one of the children until the heap property is restored



#### **Traversing down the tree**

Value	6	10	12	40	32	43	47	45	41	
Index	0	1	2	3	4	5	6	7	8	



For a node at index i, what is the index of the left and right children?

- A. (2\*i, 2\*i+1)
- B. (2\*i+1, 2\*i+2)
- C. (log(i), log(i)+1)
- **D.** None of the above

#### **Next lecture**

- More on STL implementation of heaps (priority queues)
- Queues