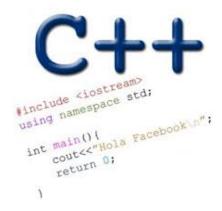
## BINARY SEARCH TREES

Problem Solving with Computers-II



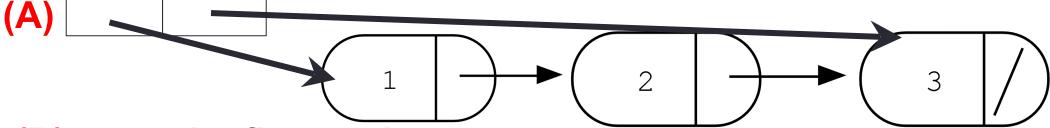
# Concept Question

```
LinkedList::~LinkedList(){
   delete head;
}
```

```
class Node {
    public:
        int info;
        Node *next;
};
```

Which of the following objects are deleted when the destructor of Linked-list is called?

head tail



(B): only the first node

(C): A and B

(D): All the nodes of the linked list

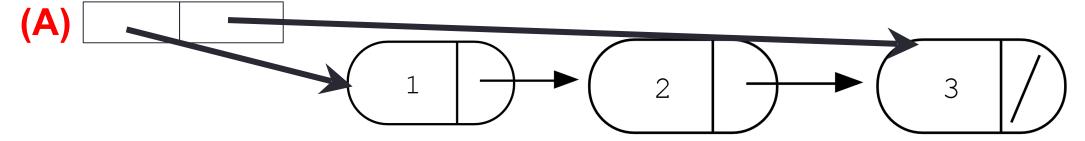
(E): A and D

## Concept question

```
LinkedList::~LinkedList(){
    delete head;
}
Node::~Node(){
    delete next;
}
```

Which of the following objects are deleted when the destructor of Linked-list is called?

head tail



(B): All the nodes in the linked-list

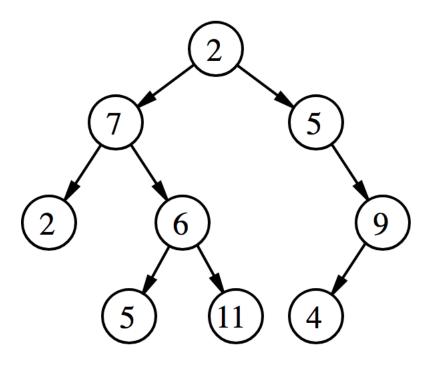
(C): A and B

(D): Program crashes with a segmentation fault

(E): None of the above

```
LinkedList::~LinkedList(){
   delete head;
}
head tail
Node::~Node(){
   delete next;
}
```

#### Trees



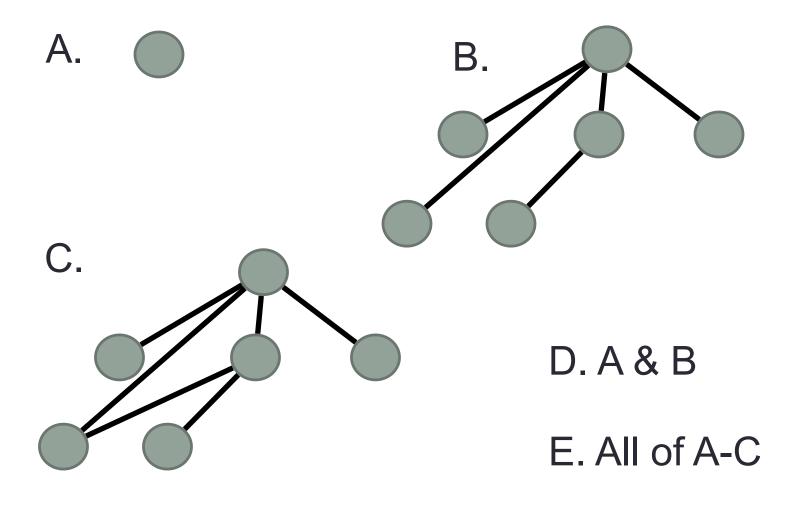
A tree has following general properties:

- One node is distinguished as a **root**;
- Every node (exclude a root) is connected by a directed edge *from* exactly one other node;

A direction is: *parent -> children* 

• Leaf node: Node that has no children

## Which of the following is/are a tree?



## Binary Search Trees

What are the operations supported?

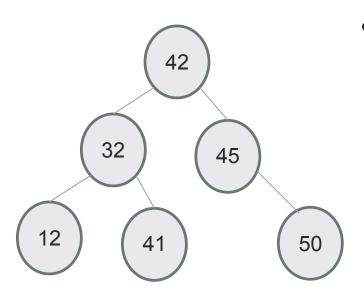
What are the running times of these operations?

How do you implement the BST i.e. operations supported by it?

### Operations supported by Sorted arrays and Binary Search Trees (BST)

Operations		
Min		
Max		
Successor		
Predecessor		
Search		
Insert		
Delete		
Print elements in order		

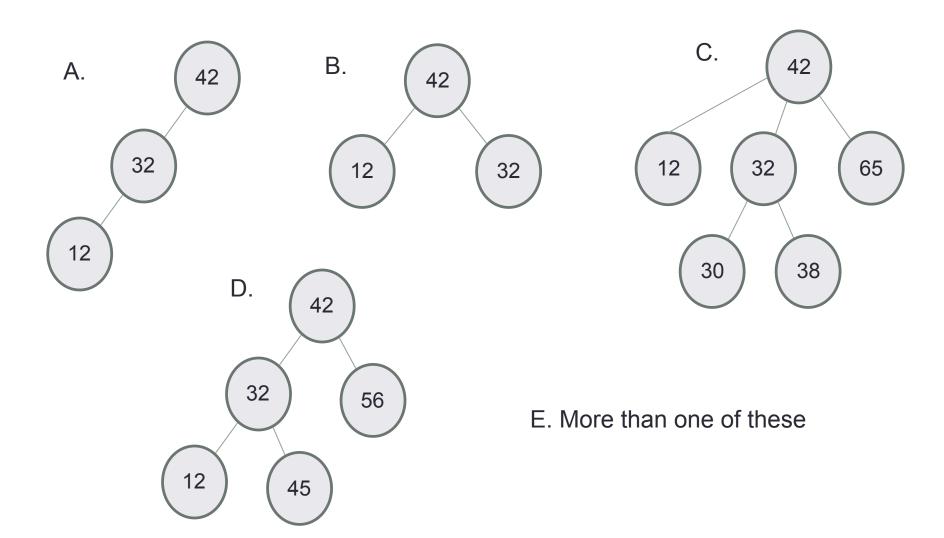
### Binary Search Tree – What is it?



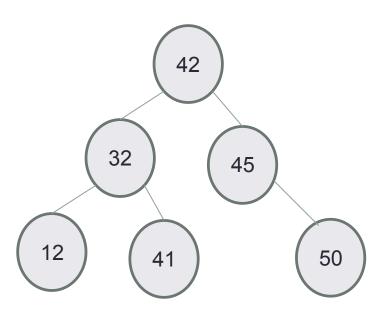
- Each node:
  - stores a key (k)
  - has a pointer to left child, right child and parent (optional)
  - Satisfies the Search Tree Property

For any node, Keys in node's left subtree <= Node's key Node's key < Keys in node's right subtree

## Which of the following is/are a binary search tree?



### BSTs allow efficient search!



- Start at the root;
- Trace down a path by comparing **k** with the key of the current node x:
  - If the keys are equal: we have found the key
  - If  $\mathbf{k} < \text{key}[\mathbf{x}]$  search in the left subtree of  $\mathbf{x}$
  - If k > key[x] search in the right subtree of x

### A node in a BST

```
class BSTNode {
public:
 BSTNode* left;
 BSTNode* right;
 BSTNode* parent;
  int const data;
 BSTNode (const int & d) : data(d) {
    left = right = parent = 0;
```

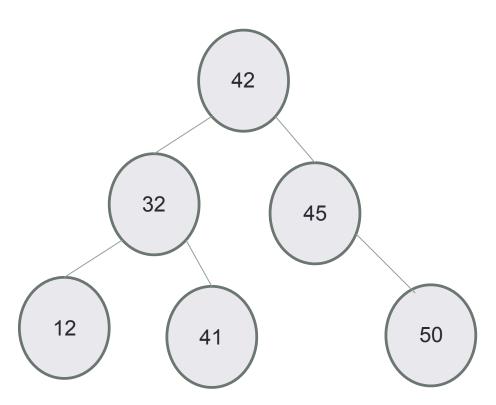
### Define the BSTADT

<b>Operations</b>	
Search	42
Insert	
Min	
Max	( 32 ) ( 45
Successor	
Predecessor	
Delete	$ \left(\begin{array}{c} 12 \right) \left(\begin{array}{c} 41 \right) \left(\begin{array}{c} 50 \right) $
Print elements in order	

## Traversing down the tree

• Suppose n is a pointer to the root. What is the output of the following code:

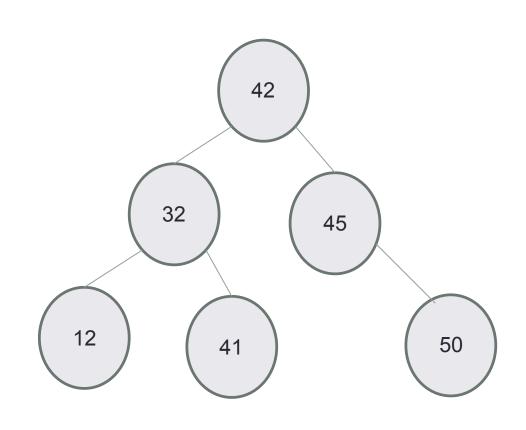
```
n = n->left;
n = n->right;
cout<<n->data<<endl;
 A. 42
 B. 32
 C. 12
 D. 41
 E. Segfault
```



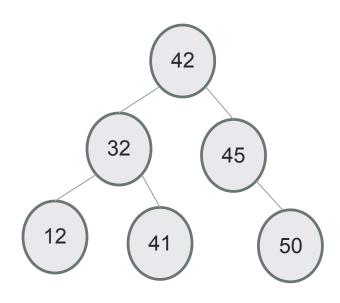
## Traversing up the tree

- Suppose n is a pointer to the node with value 50.
- What is the output of the following code:

```
n = n->parent;
  = n->parent;
n = n->left;
cout<<n->data<<endl;
 A. 42
 B. 32
 C. 12
 D. 45
 E. Segfault
```



### Insert

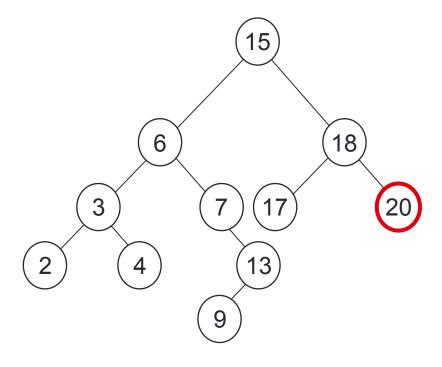


- Insert 40
- Search for the key
- Insert at the spot you expected to find it

#### Max

**Goal**: find the maximum key value in a BST Following right child pointers from the root, until a leaf node is encountered. The least node has the max value

Alg: int BST::max()



Maximum = 20

#### Min

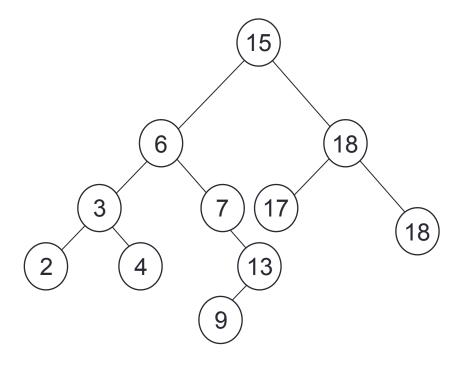
Goal: find the minimum key value in a BST

Start at the root.

Follow \_\_\_\_ child pointers from the root, until a leaf node is encountered

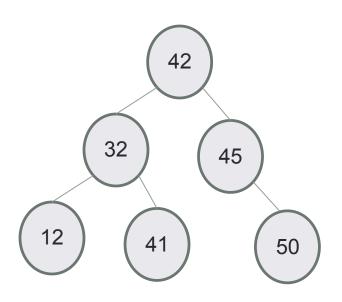
Leaf node has the min key value

Alg: int BST::min()



Min = ?

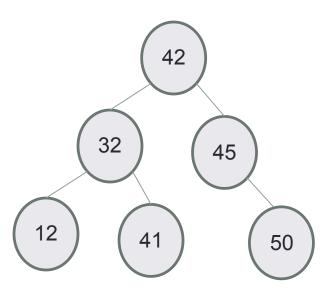
### In order traversal: print elements in sorted order



Algorithm Inorder(tree)

- 1. Traverse the left subtree, i.e., call Inorder(left-subtree)
- 2. Visit the root.
- 3. Traverse the right subtree, i.e., call Inorder(right-subtree)

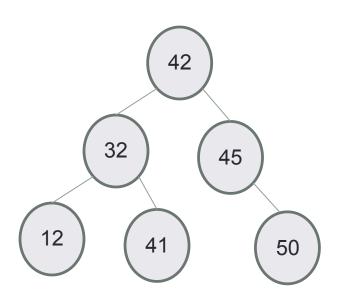
### Pre-order traversal: nice way to linearize your tree!



Algorithm Preorder(tree)

- 1. Visit the root.
- 2. Traverse the left subtree, i.e., call Preorder(left-subtree)
- 3. Traverse the right subtree, i.e., call Preorder(right-subtree)

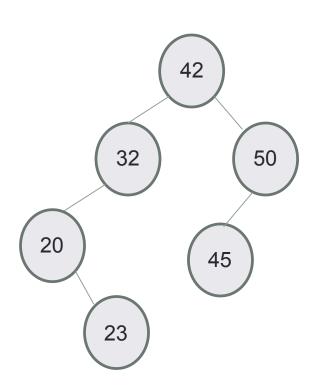
### Post-order traversal: use in recursive destructors!



Algorithm Postorder(tree)

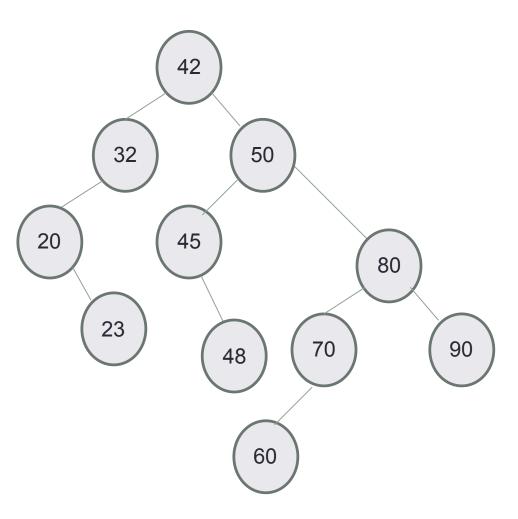
- 1. Traverse the left subtree, i.e., call Postorder(left-subtree)
- 2. Traverse the right subtree, i.e., call Postorder(right-subtree)
- 3. Visit the root.

### Predecessor: Next smallest element



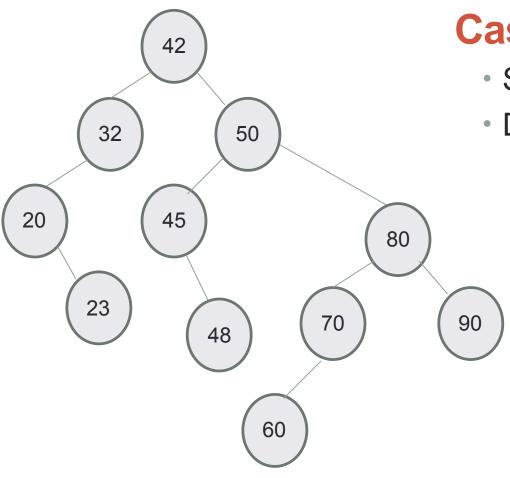
- What is the predecessor of 32?
- What is the predecessor of 45?

# Successor: Next largest element



- What is the successor of 45?
- What is the successor of 50?
- What is the successor of 60?

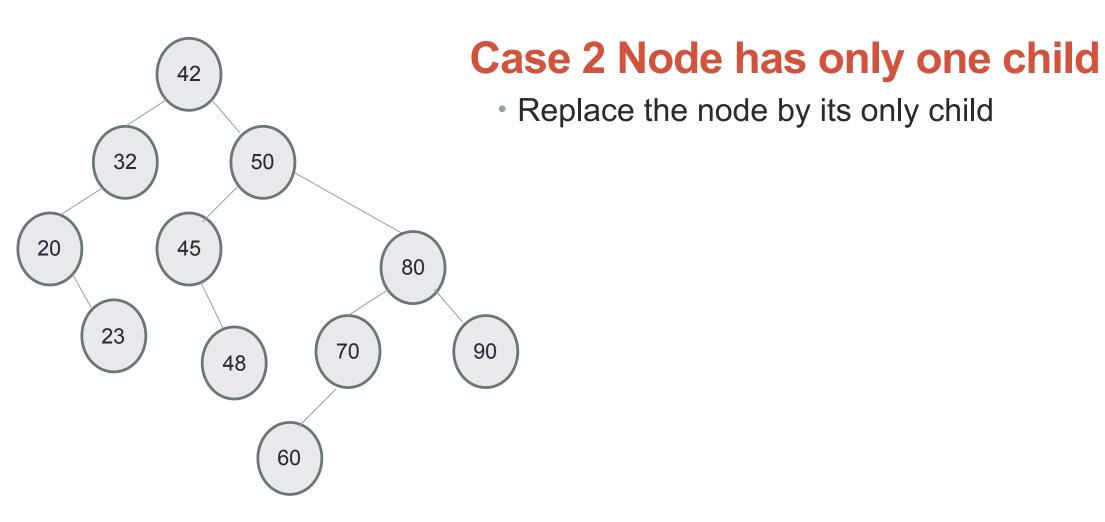
### **Delete: Case 1**



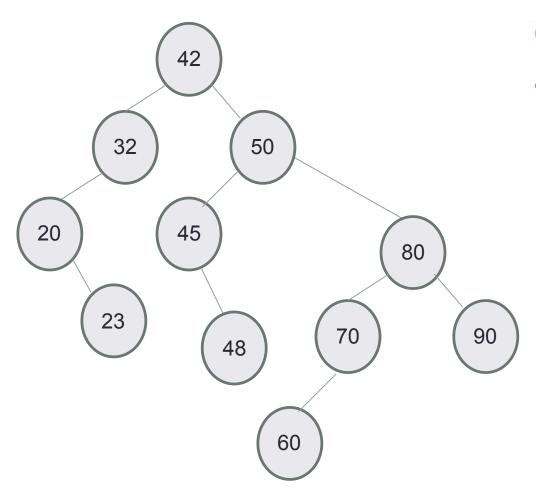
#### Case 1: Node is a leaf node

- Set parent's (left/right) child pointer to null
- Delete the node

### Delete: Case 2



### **Delete: Case 3**



#### Case 3 Node has two children

 Can we still replace the node by one of its children? Why or Why not?

### **Binary Search**

- Binary search. Given value and sorted array a[], find index i such that a[i] = value, or report that no such index exists.
- Invariant. Algorithm maintains a [lo] ≤ value ≤ a [hi].
- Ex. Binary search for 33.

