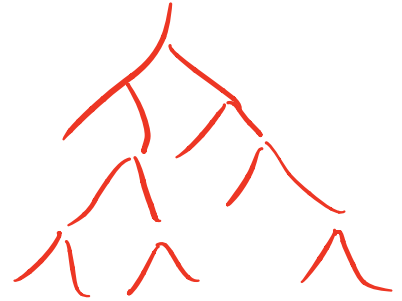


# BINARY SEARCH TREES



---

Problem Solving with Computers-II

C++

```
#include <iostream>
using namespace std;

int main(){
    cout<<"Hola Facebook!";
    return 0;
}
```

# Concept Question

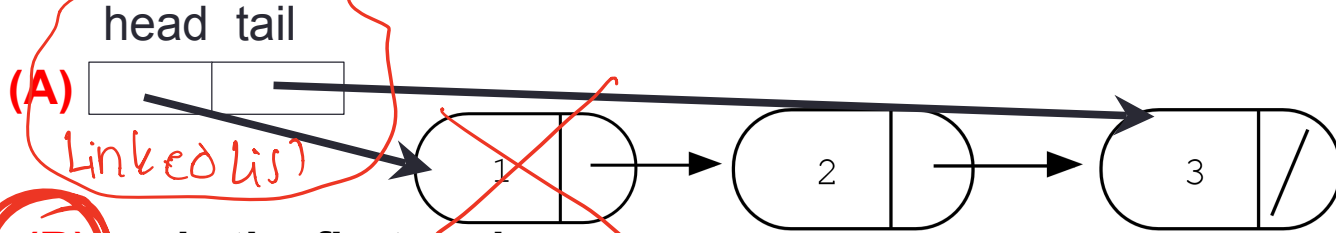
```
LinkedList::~~LinkedList(){
```

```
    delete head;
```

*private:*

```
class Node {  
    public:  
        int info;  
        Node *next;  
};
```

Which of the following objects are deleted when the destructor of Linked-list is called?



**(B): only the first node**

**(C): A and B**

**(D): All the nodes of the linked list**

**(E): A and D**

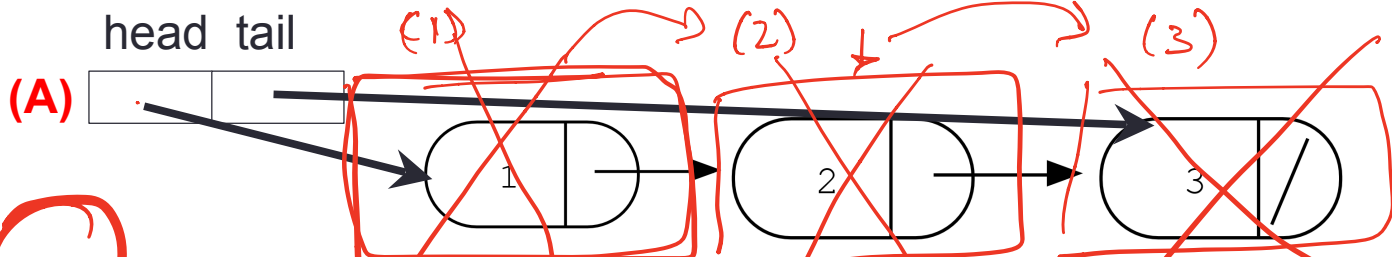
# Concept question

```
LinkedList::~~LinkedList(){  
    delete head;  
}
```

```
Node::~~Node(){  
    delete next;  
}
```

The last call to ~~destr~~ reaches the base case: delete 0;

Which of the following objects are deleted when the destructor of Linked-list is called?



**(B): All the nodes in the linked-list**

**(C): A and B**

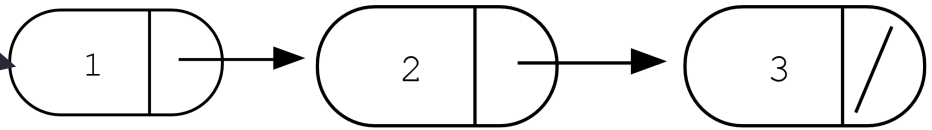
**(D): Program crashes with a segmentation fault**

**(E): None of the above**

delete 0;  
Notes: Deleting a null pointer does not result in a segfault

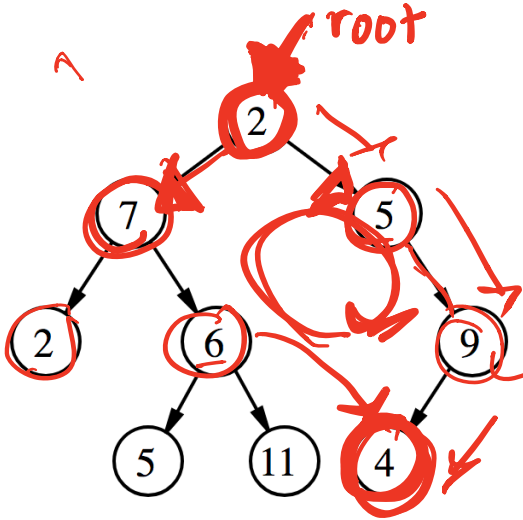
```
LinkedList::~~LinkedList(){
    delete head;
}
```

```
Node::~~Node(){
    delete next;
}
```



Binary tree: At most 2-children

# Trees Hierarchical



A tree has following general properties:

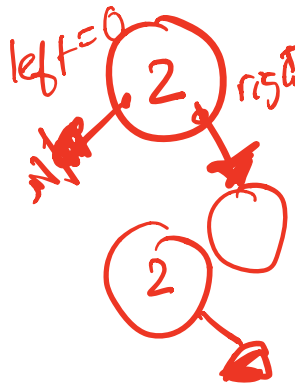
- One node is distinguished as a **root**;
- Every node (exclude a root) is connected by a directed edge *from* exactly one other node;

(key 2)

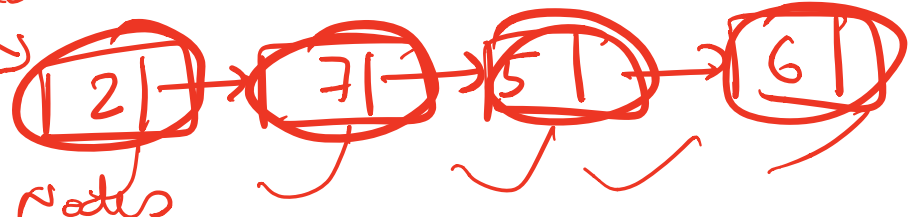
A direction is: *parent -> children*

- *Leaf node: Node that has no children*

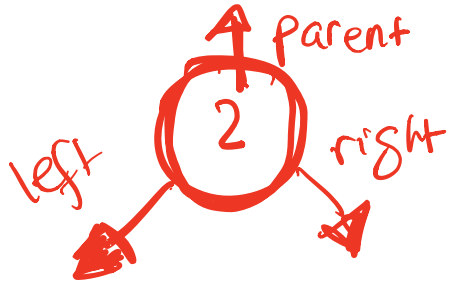
key (data)  
parent-child



head



Linear structure.



```
class Node {  
    public:  
        int data;  
        Node * parent;  
        Node * left;  
        Node * right;  
};
```

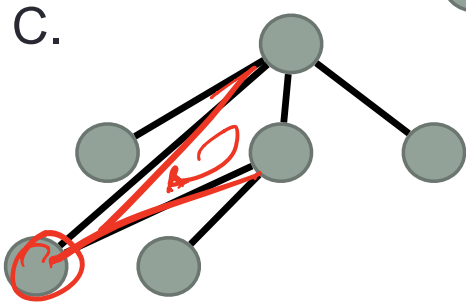
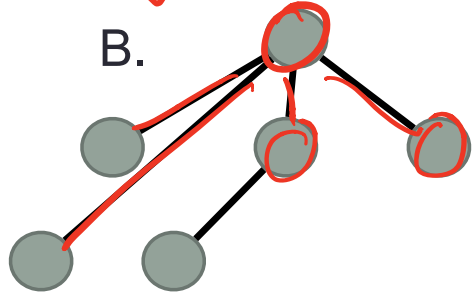
Which of the following is/are a tree?

Node  $\neq$

root = 0  
empty tree



Single node tree



D. A & B

E. All of A-C

# Binary Search Trees

(Variant)

Red-Black  
AVL

- What are the operations supported?

insert, delete, search, min, max, Vanilla!

- What are the running times of these operations?

Build intuition → Complexity

- How do you implement the BST i.e. operations supported by it?

~~\_\_\_\_\_~~

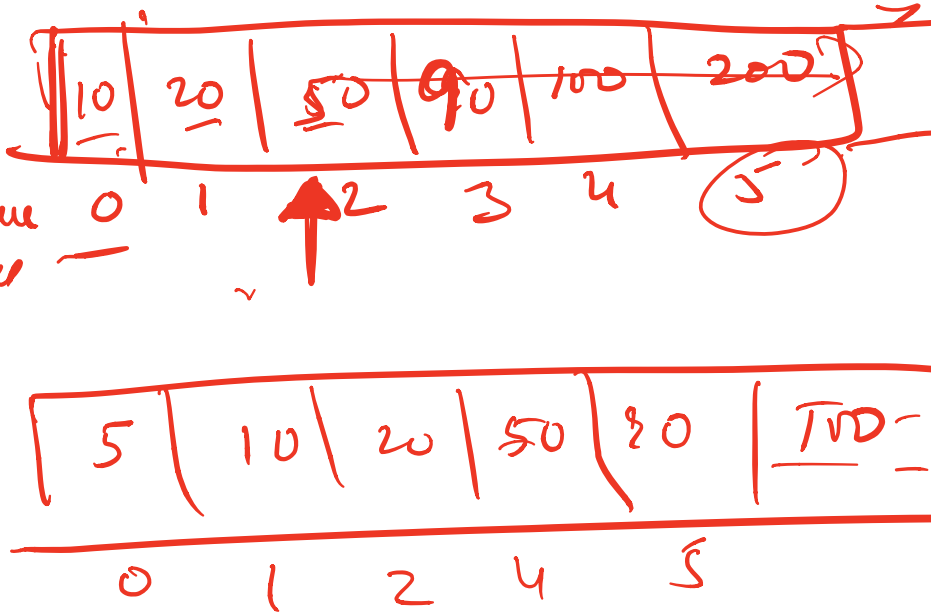


# Operations supported by Sorted arrays and Binary Search Trees (BST)

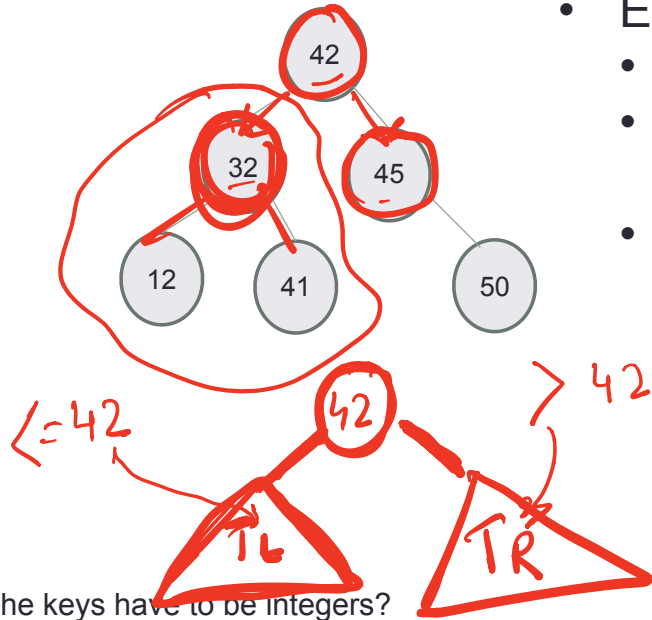


Search property

Operations	
Min	
Max	
Successor	next largest value
Predecessor	next smaller value
Search	Binary Search
Insert	
Delete	Slow (fast)
Print elements in order	



# Binary Search Tree – What is it?



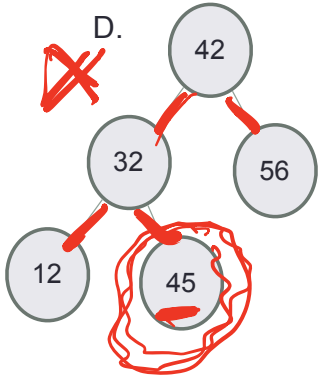
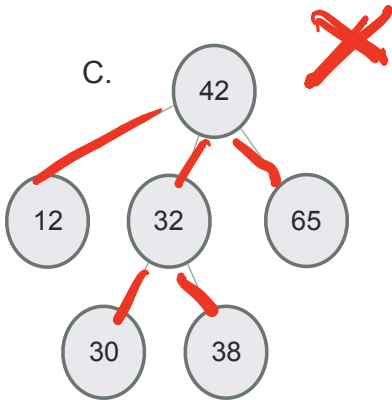
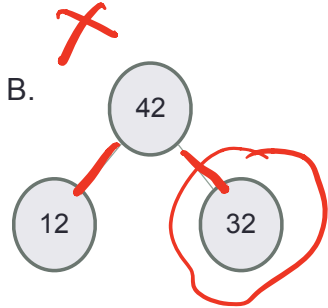
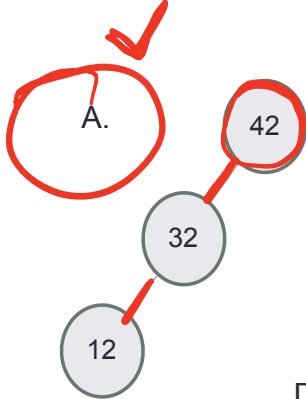
- Each node:
  - stores a key ( $k$ )
  - has a pointer to left child, right child and parent (optional)
  - Satisfies the **Search Tree Property**

For any node,  
 Keys in node's left subtree  $\leq$  Node's key  
 Node's key  $<$  Keys in node's right subtree

$$\text{keys}(T_L) \leq \text{key}(x) < \text{key}(T_R)$$

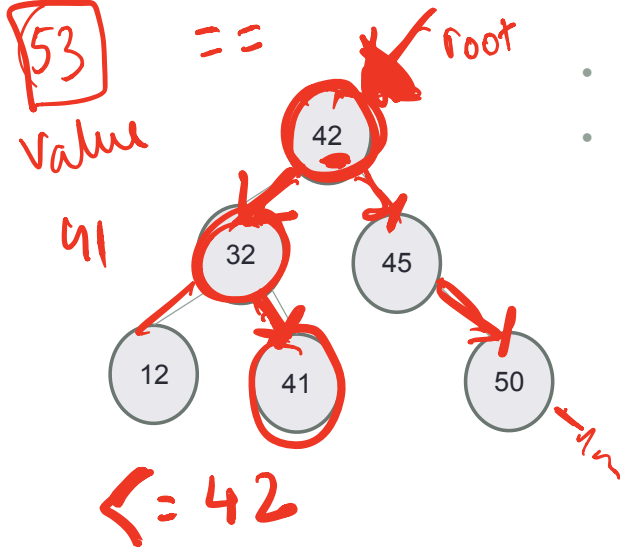
(BST)

Which of the following is/are a binary search tree?



E. More than one of these

# BSTs allow efficient search!



- Start at the root;
- Trace down a path by comparing  $k$  with the key of the current node  $x$ :
  - If the keys are equal: we have found the key
  - If  $k < \text{key}[x]$  search in the left subtree of  $x$
  - If  $k > \text{key}[x]$  search in the right subtree of  $x$



Search for 41, then search for 53

## A node in a BST

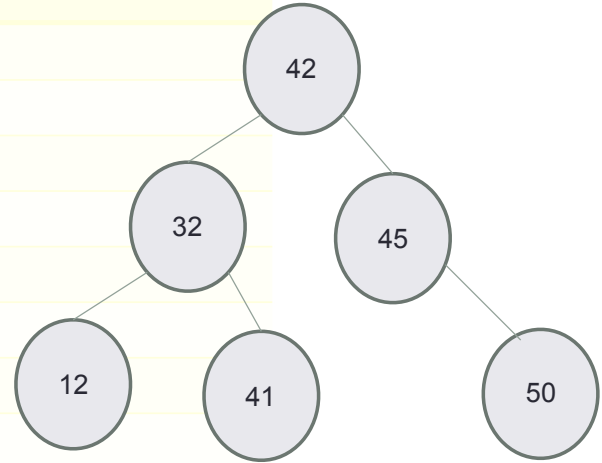
```
class BSTNode {  
  
public:  
    BSTNode* left;  
    BSTNode* right;  
    BSTNode* parent;  
    int const data;  
  
    BSTNode( const int & d ) : data(d) {  
        left = right = parent = 0;  
    }  
};
```

*initializer list*

# Define the BST ADT

Abstract Data Type

<b>Operations</b>
Search
Insert
Min
Max
Successor
Predecessor
Delete
Print elements in order



# Traversing down the tree

- Suppose `n` is a pointer to the root. What is the output of the following code:

```
n = n->left;
```

```
n = n->right;
```

```
cout<<n->data<<endl;
```

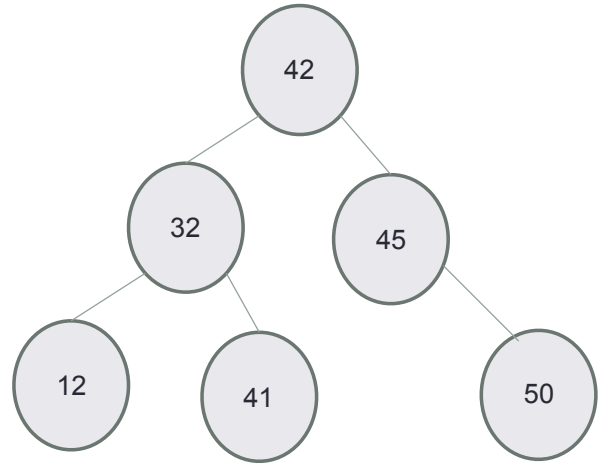
A. 42

B. 32

C. 12

D. 41

E. Segfault



# Traversing up the tree

- Suppose `n` is a pointer to the node with value 50.
- What is the output of the following code:

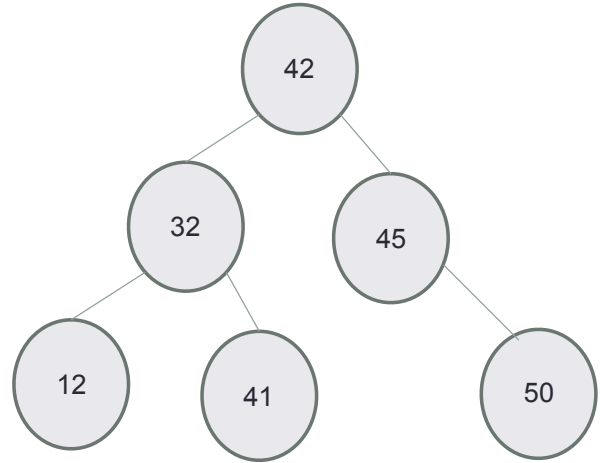
```
n = n->parent;
```

```
n = n->parent;
```

```
n = n->left;
```

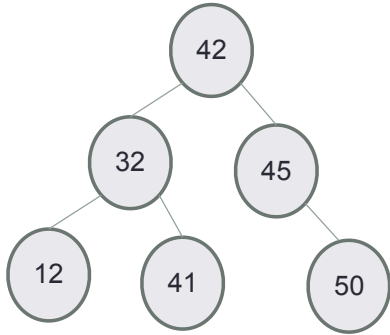
```
cout<<n->data<<endl;
```

- A. 42
- B. 32
- C. 12
- D. 45
- E. Segfault





# Insert



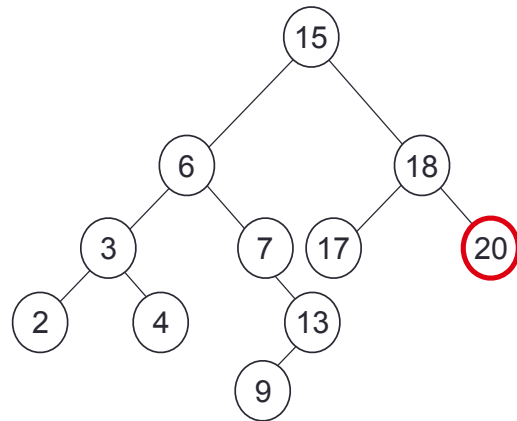
- Insert 40
- Search for the key
- Insert at the spot you expected to find it

# Max

**Goal:** find the maximum key value in a BST

Following right child pointers from the root, until a leaf node is encountered. The least node has the max value

**Alg:** `int BST::max()`



**Maximum = 20**

# Min

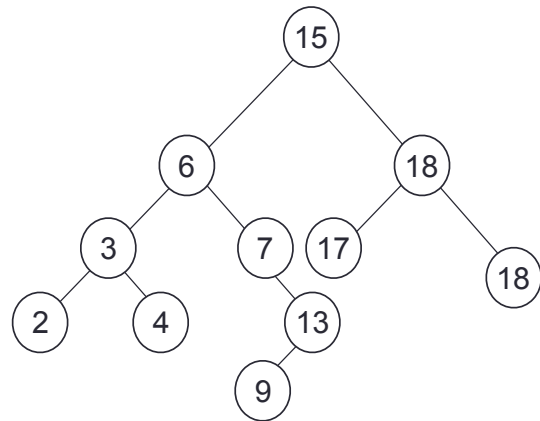
**Goal:** find the minimum key value in a BST

Start at the root.

Follow \_\_\_\_\_ child pointers from the root, until a leaf node is encountered

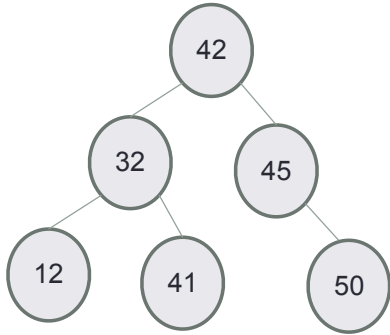
Leaf node has the min key value

**Alg:** `int BST::min()`



Min = ?

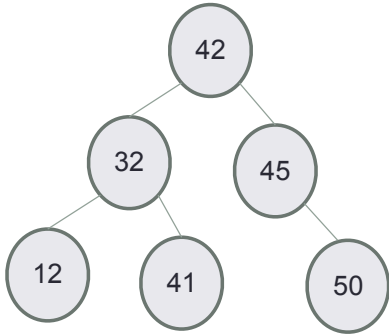
# In order traversal: print elements in sorted order



Algorithm Inorder(tree)

1. Traverse the left subtree, i.e., call Inorder(left-subtree)
2. Visit the root.
3. Traverse the right subtree, i.e., call Inorder(right-subtree)

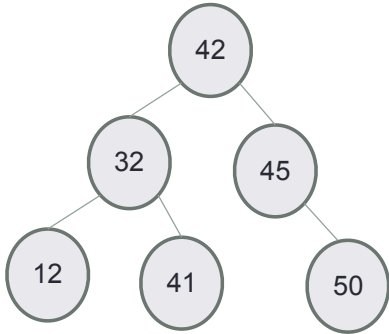
# Pre-order traversal: nice way to linearize your tree!



Algorithm Preorder(tree)

1. Visit the root.
2. Traverse the left subtree, i.e., call Preorder(left-subtree)
3. Traverse the right subtree, i.e., call Preorder(right-subtree)

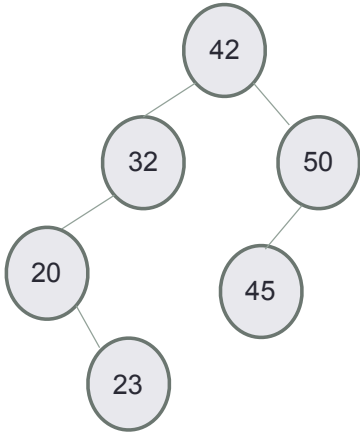
# Post-order traversal: use in recursive destructors!



Algorithm Postorder(tree)

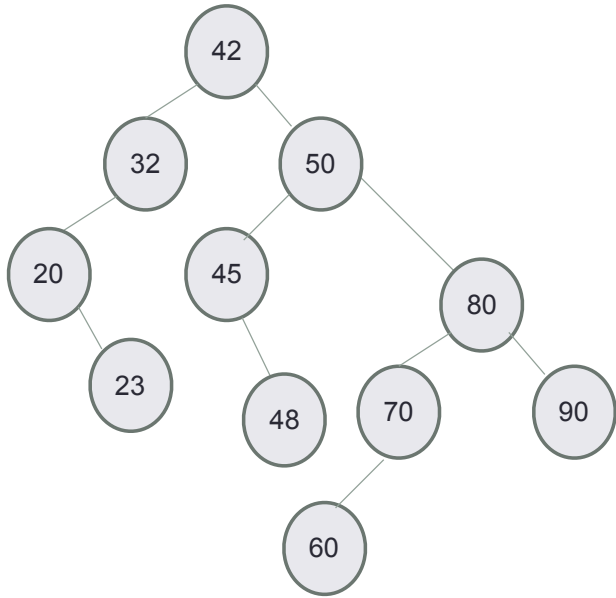
1. Traverse the left subtree, i.e., call Postorder(left-subtree)
2. Traverse the right subtree, i.e., call Postorder(right-subtree)
3. Visit the root.

# Predecessor: Next smallest element



- What is the predecessor of 32?
- What is the predecessor of 45?

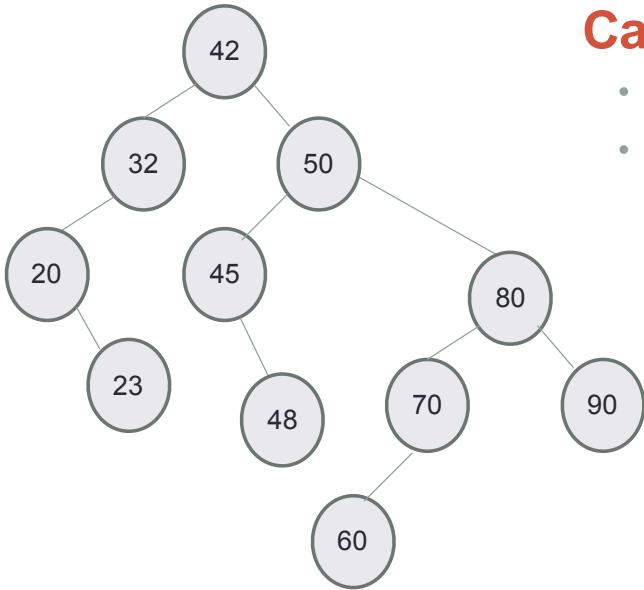
# Successor: Next largest element



- What is the successor of 45?
- What is the successor of 50?
- What is the successor of 60?



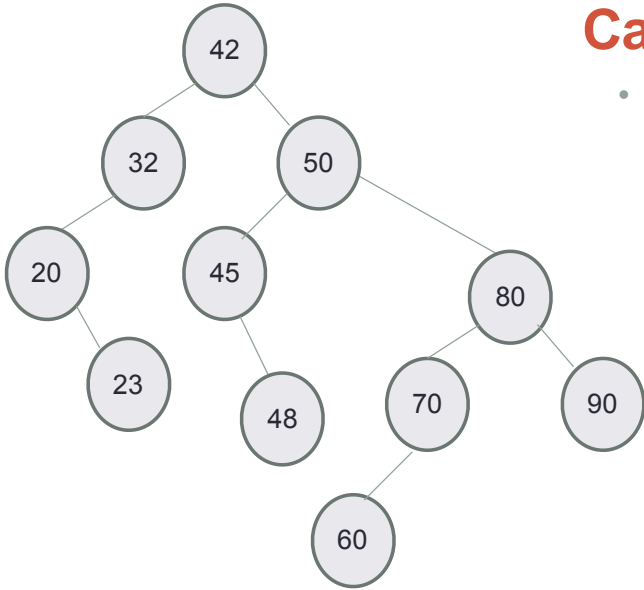
# Delete: Case 1



## Case 1: Node is a leaf node

- Set parent's (left/right) child pointer to null
- Delete the node

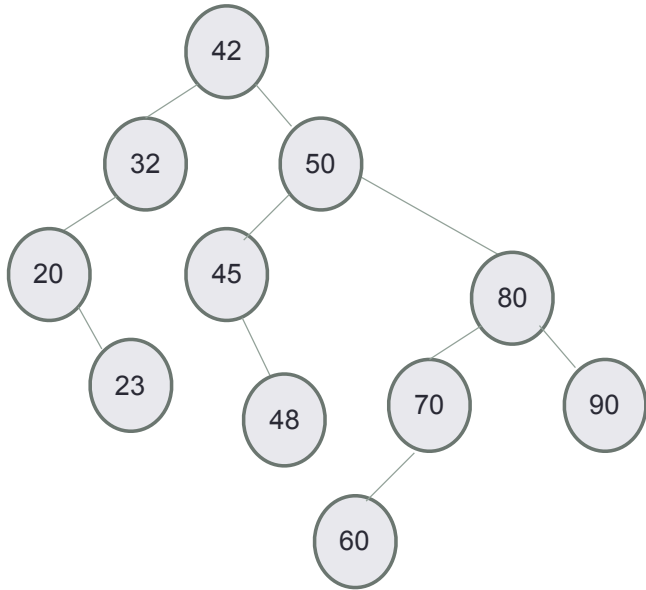
# Delete: Case 2



## Case 2 Node has only one child

- Replace the node by its only child

# Delete: Case 3



## Case 3 Node has two children

- Can we still replace the node by one of its children? Why or Why not?

