# QUEUES AND PRIORITY QUEUES

**Problem Solving with Computers-II** 





# How is PA03 going?

- A. Done
- B. On track to finish
- C. Having trouble with the checkpoint (design)
- D. Just started
- E. Haven't started

# Evaluating a fully parenthesized infix expression

Characters read so far (shaded): (((6 + 9) / 3) \* (6 - 4))

Numbers







Lab 05, part2 : Evaluating post fix expressions using a single stack Postfix: 7 3 5 \* + 4 2 / -Infix: (7 + (3 \* 5)) - (4 / 2)

4

### The Queue Operations

- A queue is like a line of people waiting for a bank teller.
- The queue has a **front** and a **rear**.



# The Queue Operations

New people must enter the queue at the rear. The C++ queue class calls this a <u>push</u>, although it is usually called an <u>enqueue</u> operation.





### The Queue Operations

 When an item is taken from the queue, it always comes from the front. The C++ queue calls this a <u>pop</u>, although it is usually called a <u>dequeue</u> operation.



# The Queue Class

- The C++ standard template library has a queue template class.
- The template parameter is the type of the items that can be put in the queue.

```
template <class Item>
class queue<Item>
{
public:
    queue( );
    void push(const Item& entry);
    void pop( );
    bool empty( ) const;
    Item front( ) const;
    ...
```

### Queue via stacks

Implement a MyQueue class which implements a queue using two stacks





# Priority Queues or Heaps

Min-Heaps Max-Heap

BST

- Insert :
- Min:
- Delete Min:
- Max
- Delete Max

Choose heap if you are doing repeated insert/delete/(min OR max) operations

#### **Applications:**

- Efficient sort
- Finding the median of a sequence of numbers
- Compression codes

### Heaps as binary trees

- Rooted binary tree that is as complete as possible
- In a min-Heap, each node satisfies the following heap property: key(x)<= key(children of x)</li>



### Heaps as binary trees

- Rooted binary tree that is as complete as possible
- In a max-Heap, each node satisfies the following heap property: key(x)>= key(children of x)



# Identifying heaps

Starting with the following min Heap which of the following operations will result in something that is NOT a min Heap

- A. Swap the nodes 40 and 32
- B. Swap the nodes 32 and 43
- C. Swap the nodes 43 and 40
- D. Insert 50 as the left child of 45 E. C&D



### Structure: Complete binary tree

A heap is a complete binary tree: Each level is as full as possible. Nodes on the bottom level are as far left as possible



## Insert 32 into a heap

- Insert key(x) in the first open slot at the last level of tree (going from left to right)
- If the heap property is not violated Done
- Else: while(key(parent(x))>key(x)) swap the key(x) with key(parent(x))



### Insert 50, then 35, then 8



### **Delete min**

- Replace the root with the rightmost node at the last level
- "Bubble down"- swap node with one of the children until the heap property is restored



### Next lecture

• Under the hood of priority queues