

FINAL PRACTICE!

Problem Solving with Computers-II

Link to class handout:

<https://bit.ly/one-problem-to-rule-them-all>

The image shows the C++ logo in blue, with the text 'C++' in a bold, sans-serif font. Below the logo is a snippet of C++ code in a monospaced font, with syntax highlighting: '#include <iostream>', 'using namespace std;', 'int main(){', 'cout<<"Hola Facebook!n";', 'return 0;', and '}'.

```
#include <iostream>
using namespace std;

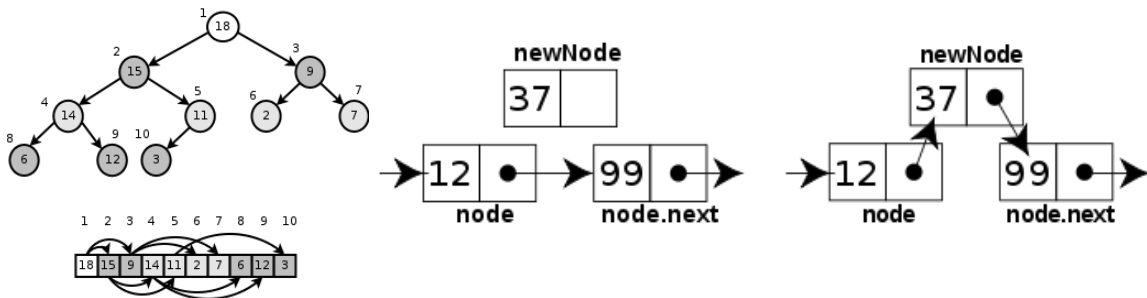
int main(){
    cout<<"Hola Facebook!n";
    return 0;
}
```

I can deal with pressure, and deadlines.



Resources for the Final Exam

- Leet code problems categorized to help you focus on the most relevant ones.
 - Those marked as “Optional challenge” won’t be on the final
- Code from lectures: <https://github.com/ucsb-cs24-w26/cs24-w26-lectures>
- Last year’s final exam posted on Canvas for Practice (under module 10)
- Tool to visualize data structures: <https://visualgo.net/>
- If you missed your Mock interview, complete during section tomorrow.



INSERTION-SORT(A)	<i>cost</i>	<i>times</i>
1 for $j = 2$ to $A.length$	c_1	n
2 $key = A[j]$	c_2	$n - 1$
3 // Insert $A[j]$ into the sorted sequence $A[1..j - 1]$.	0	$n - 1$
4 $i = j - 1$	c_4	$n - 1$
5 while $i > 0$ and $A[i] > key$	c_5	$\sum_{j=2}^n t_j$
6 $A[i + 1] = A[i]$	c_6	$\sum_{j=2}^n (t_j - 1)$
7 $i = i - 1$	c_7	$\sum_{j=2}^n (t_j - 1)$
8 $A[i + 1] = key$	c_8	$n - 1$

Data Structures and C++

Complexity Analysis

One problem to rule them all!

Problem (LP04): Longest Consecutive Sequence

Description: Given an unsorted array of integers `nums`, return the length of the longest consecutive elements sequence.

Try solving it in as many different ways as you can, then pick the best!

Ex: `nums = [100, 4, 200, 1, 3, 2]`

output = 4

`nums = [1, 0, 1, 2]`

output = 3

Approach 1 → Sort + ?

nums = [100, 4, 200, 1, 3, 2]

nums = [1, 0, 1, 2]

= [1, 2, 3, 4, 100, 200]
nums[i] + 1 == nums[i+1]?

i				
0	1	+	1	== 2 ✓ (T)
1	2	+	1	== 3 ✓ (T)
2	3	+	1	== 4 ✓ (T)
3	4	+	1	== 100 ✗ (F)
4	100	+	1	== 200 ✗ (F)

seq len = 1 max len = 1

- 2
- 3
- 4
- 1
- 1

{ ① max(max len, seq len) = 4
 { ② seq len = 1
 ↑ repeat

→ **max len = 4**

nums = [1, 0, 1, 2]
 ↓ ↓
 0, 1, 1, 2

Think through case with
duplicates.

i
0
1
2
 seq len = 1
 2
 2
 3

If there is a duplicate do nothing

if (nums[i] == nums[i+1])
 continue // do nothing

Approach 2 - Use a Data Structure that maintains order (std::set)
or partial order (priority-queue)

nums = [100, 4, 200, 1, 3, 2]

nums = [1, 0, 1, 2]

num: num s

1

2

3

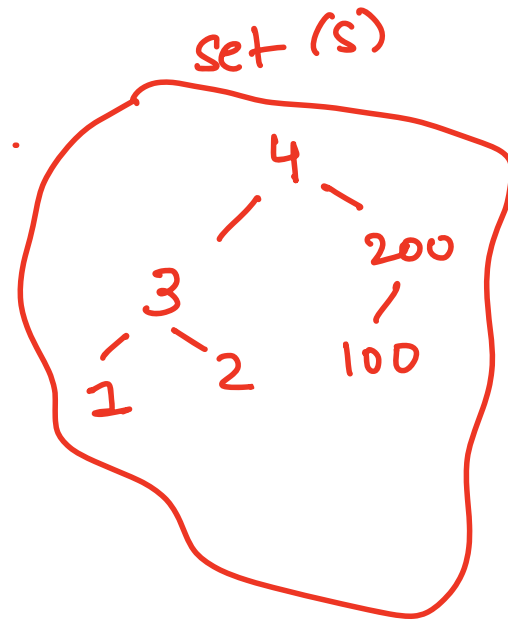
4

100

200

check if next
element is consecutive

- ① $s.count(num+1) == 1$
 - ② $s.contains(num+1)$
 - ③ Use an iterator
- } ways to find a key in a set



Approach 3 no sorting or sorted DS.

nums = [100, 4, 200, 1, 3, 2]

num: H

→ 3

→ 1

Avg. O(1) find, contains, count

H. contains (4) ✓

H. contains (5) X

H. contains (2) ✓

H. " (3)

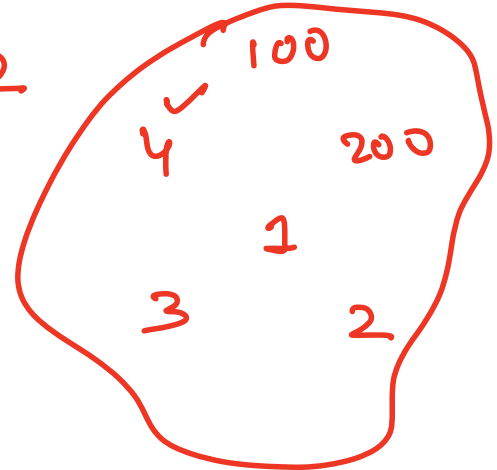
H. " (4) ✓

H. " (5) X

seq len = 2

seq len = 4

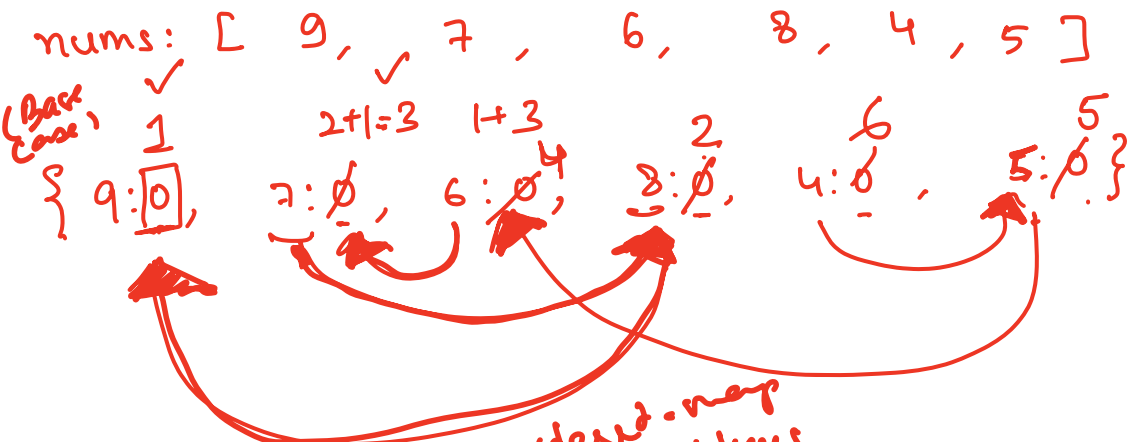
hashtable (H)



↓
You can also visualize as

{ 100, 4, 200, 1, 3, 2 }

Approach 4 ?



Pseudocode

Initialize $M[key] = 0$ for all keys

max-len = 0

for each key (k) in M:

max-len = max(max-len, lcs(M, k))

Approach: Recursively compute the lcs for each key.

Use an unordered_map to store the lcs starting at each key.

and reuse computed values.

also called memoization

Pseudo code .

```

lcs (M: unordered_map, key: int)
if key ∉ M: // key does not exist
return 0
if M[key] > 0
return M[key]
return 1+lcs(M, key+1)

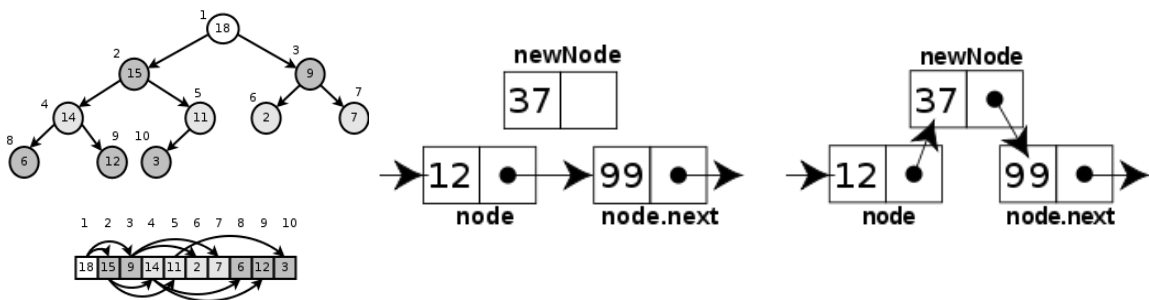
```

CS24: Learning goals

Correct coding, clear thinking, no AI shortcuts

- Design and implement **larger programs** that **run fast**
- Organize **data** in programs using **data structures**
- **Analyze** the **complexity** of your programs
- Prep for **technical interviews**

Grow confident in your problem solving skills!



INSERTION-SORT(A)

```

1 for  $j = 2$  to  $A.length$ 
2    $key = A[j]$ 
3   // Insert  $A[j]$  into the sorted
   sequence  $A[1..j-1]$ .
4    $i = j - 1$ 
5   while  $i > 0$  and  $A[i] > key$ 
6      $A[i + 1] = A[i]$ 
7      $i = i - 1$ 
8    $A[i + 1] = key$ 

```

cost *times*

c_1 n

c_2 $n - 1$

0 $n - 1$

c_4 $n - 1$

c_5 $\sum_{j=2}^n t_j$

c_6 $\sum_{j=2}^n (t_j - 1)$

c_7 $\sum_{j=2}^n (t_j - 1)$

c_8 $n - 1$

Data Structures and C++

Complexity Analysis

Break: Please take a moment to fill the course evaluations!



PROBLEM SOLVING II

Student-FO

<https://go.blueja.io/2Py4AgwuLkGbGaMq5b8sKg>

To access the evaluation, scan this QR code with your mobile phone.

Thank you and all the best !

