



# QUEUES

# BREADTH-FIRST TRAVERSAL

# COMPLETE BINARY TREES

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Problem Solving with Computers-II

# C++

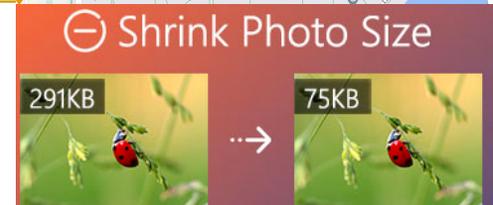
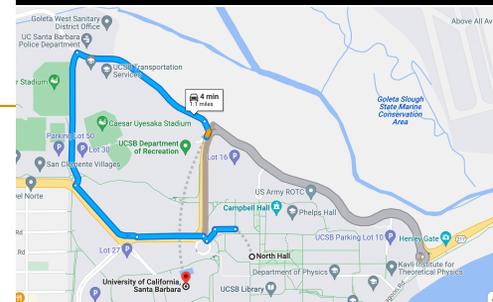
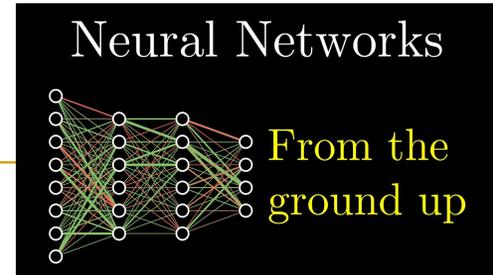
```
#include <iostream>
using namespace std;

int main(){
    cout<<"Hola Facebook!";
    return 0;
}
```

GitHub

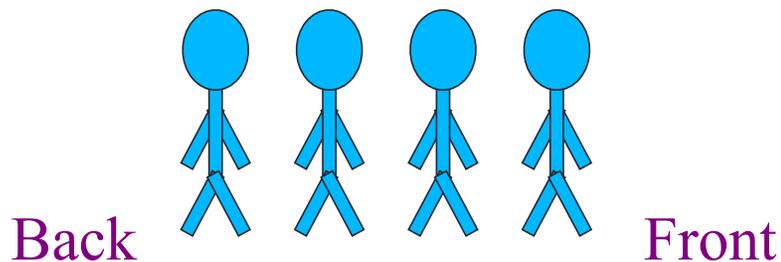


ADT	Algorithm	Real-World Application
<b>Stack</b>	Depth-First Search	Machine Learning <b>(PA03: Training and Prediction in Neural Nets)</b>
<b>Queue</b>	Breadth-First Search	
<b>Priority Queue (Lab 04)</b>	Dijkstra's Algorithm	GPS Navigation (Shortest path)
<b>Priority Queue</b>	Huffman Coding	Data Compression (ZIP, JPEG, MP3)
<b>Your choice!</b>	You design!	Querying a movie dataset <b>(PA02)</b>



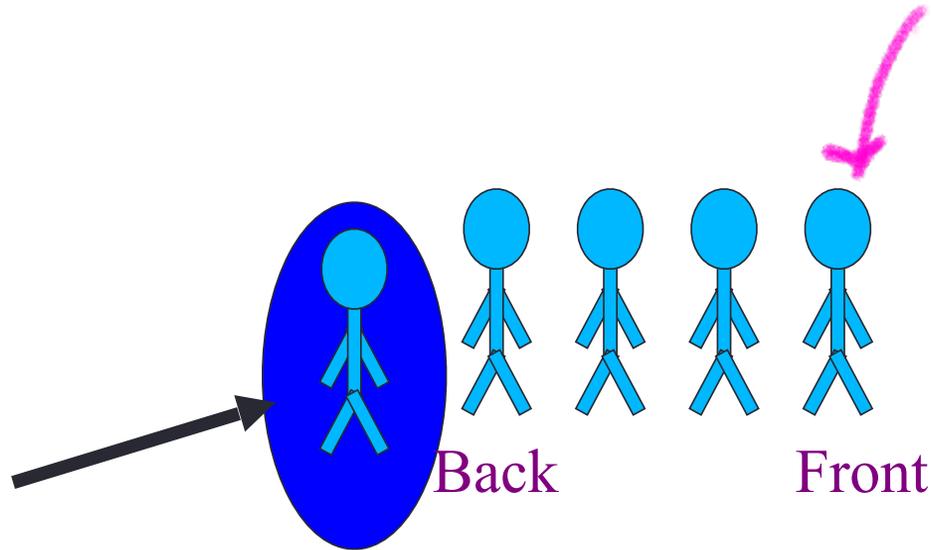
# Queue: First come First Serve

- A queue is like a queue of people waiting to be serviced
- The queue has a front and a back.



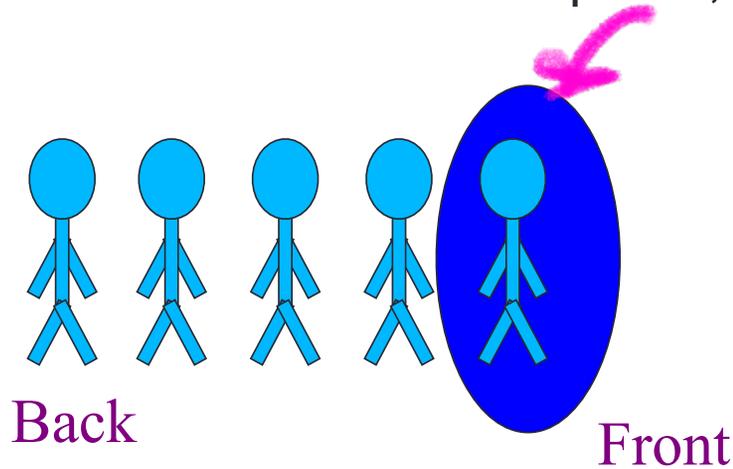
# Queue Operations: push, pop, front, back

New people must enter the queue at the back. The C++ queue class calls this a push operation.



# Queue Operations: push, pop, front, back - $O(1)$

- To check the item in the front of the queue, use **front()**
- To check the item at the back of the queue, use **back()**
- When an item is taken from the queue, it always comes from the front.
- To delete an element from the front of the queue, use **pop()**



# Queue Operations: empty(), push, pop, front, back: **O(1)**

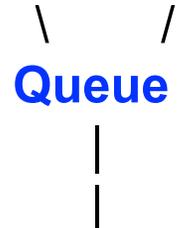
```

std::queue<int> q;
q.empty(); //true
q.push(1);
// push 2, 3, 4, 5
q.front();
q.back();
q.pop();

```

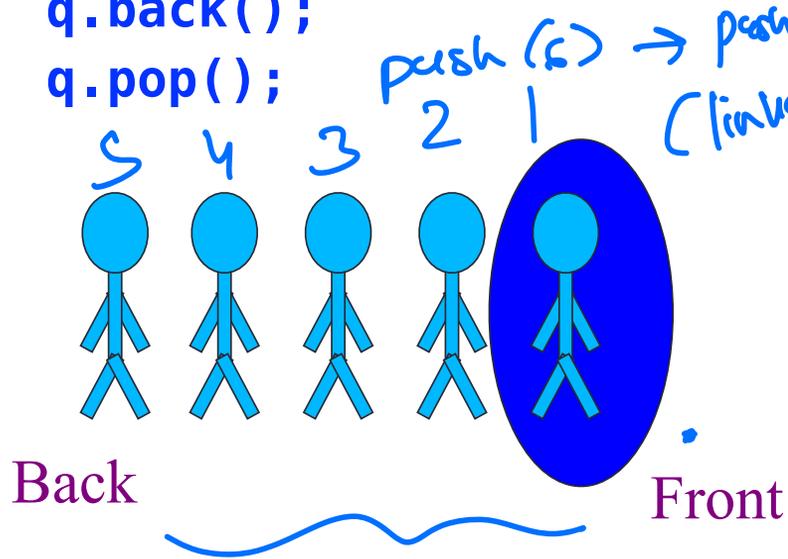
Algorithms: Breadth First Search Task Scheduling

ADT:

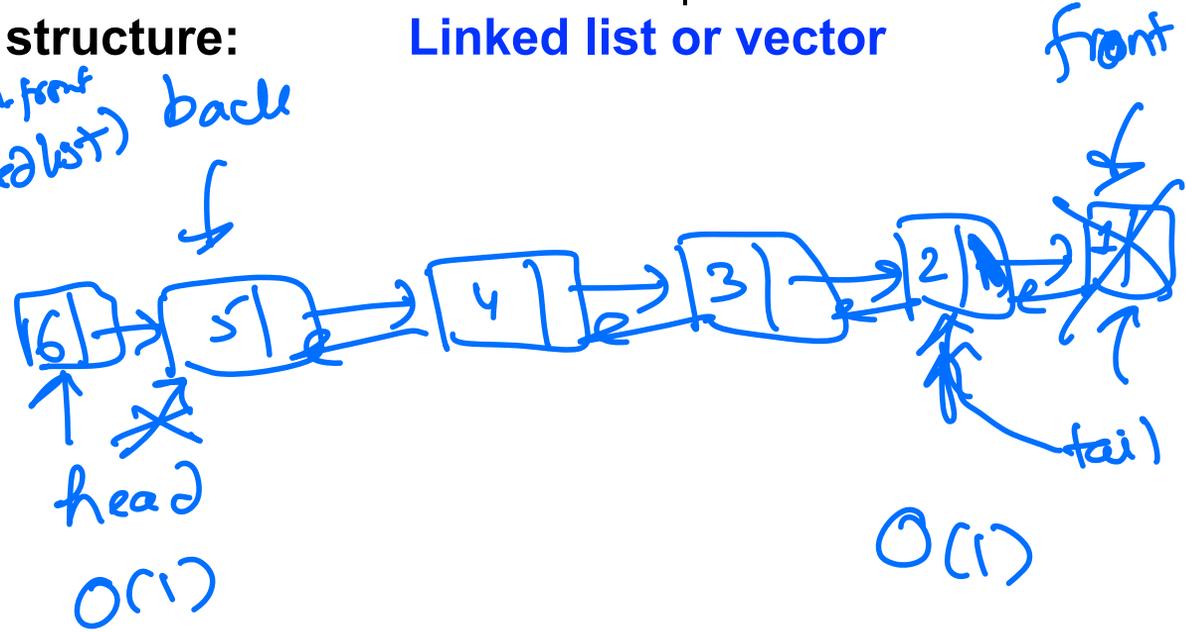


Data structure:

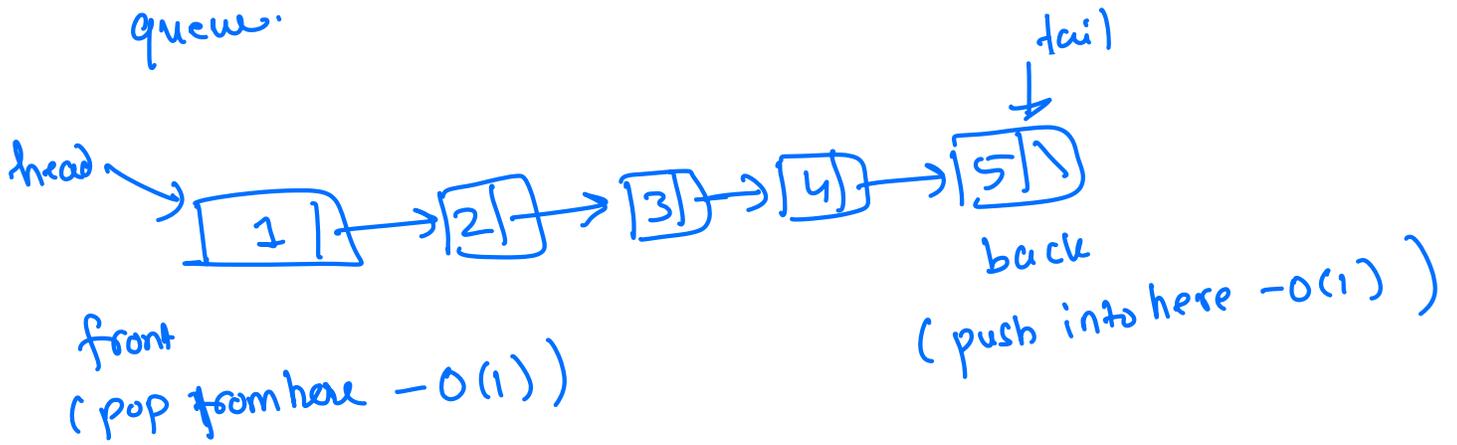
Linked list or vector

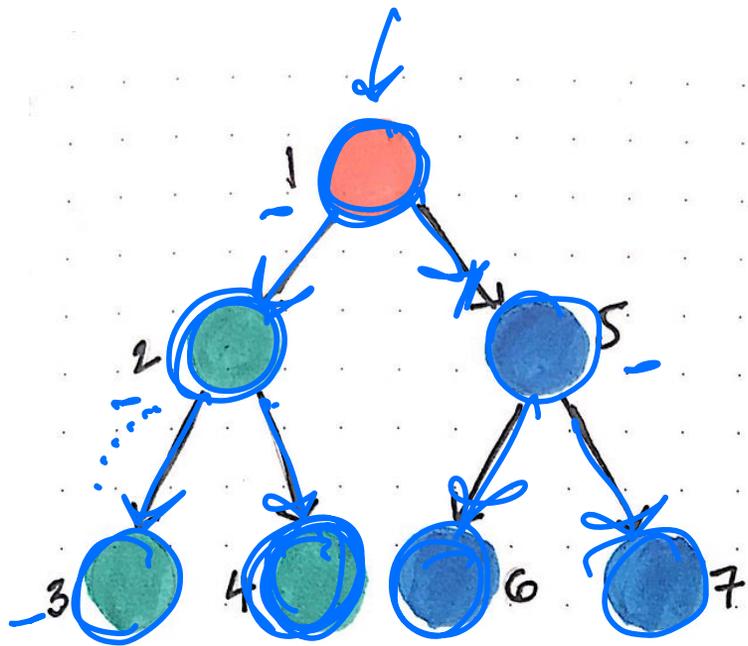


push (6) → push front (linked list) back



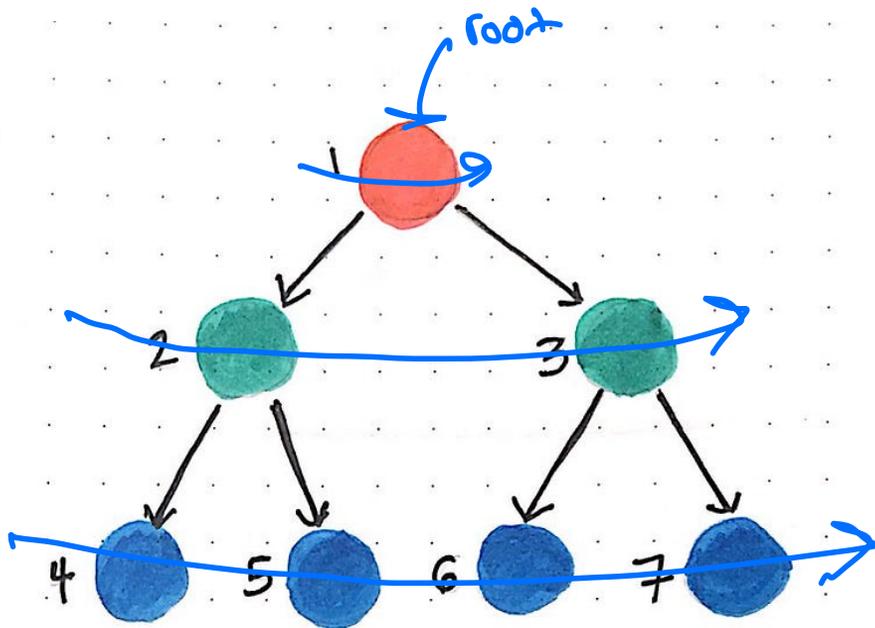
Another possible implementation for a queue is to have a singly-linked list with head of list representing the front of the queue and tail of list representing the back of the queue.





### Depth-first search

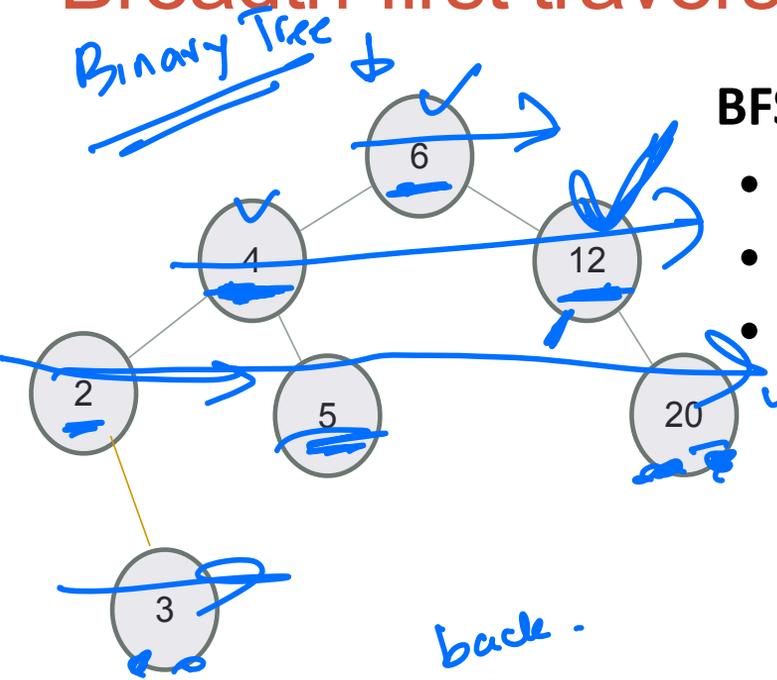
- Traverse through left subtree(s) first, then traverse through the right subtree(s).



### Breadth-first search

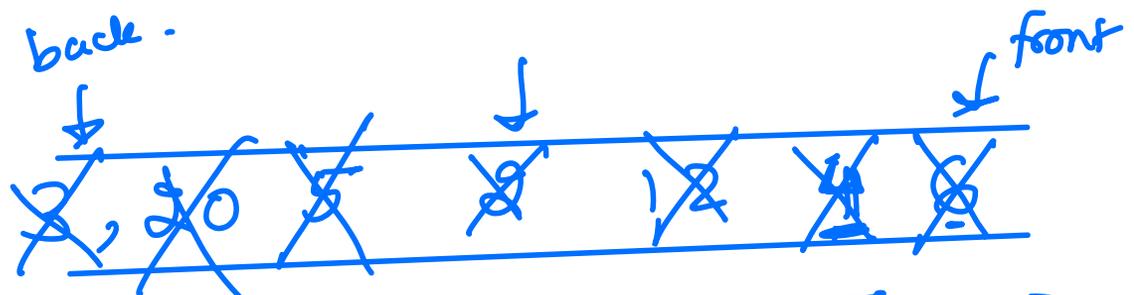
- Traverse through one level of children nodes, then traverse through the level of grandchildren nodes (and so on...).

# Breadth-first traversal/search



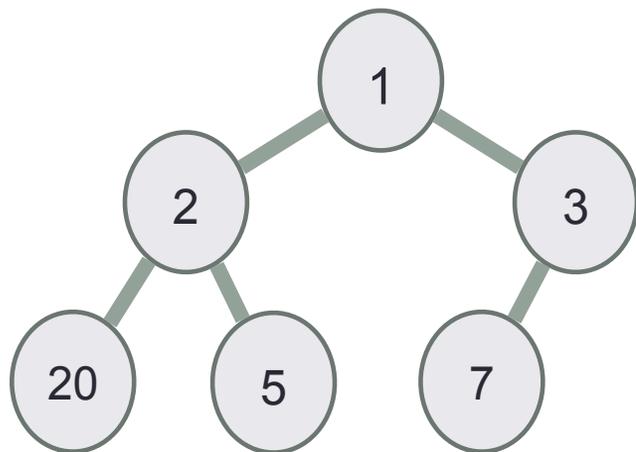
## BFS Algo:

- Create an empty queue.
- Push Insert the **root** into the **queue**.
- While queue is not empty,
  - Print the key in the front of the queue
  - Push all the children of the node into the queue.
  - Pop the front node from the queue



Output : 6, 4, 12, 2, 5, 20, 3

# Breadth-first traversal



**BFS Algo (store output in a vector: result):**

- Create an empty **queue**.
- Create an empty **vector** called **result**.
- Insert the **root** into the **queue**.
- While queue is not empty,
  - **Append the key in the front of the queue to result**
  - Push all the children of the node into the queue.
  - Pop the front node from the queue

**Activity 1:**

1. Write the **Breadth First Traversal** in your handout
2. Trace it on the given tree, showing how the queue evolve.
3. Show the output as a vector of key values?

# Connect vector with Google maps!

Applications: Machine Learning, Operating System    Image compression, Google maps

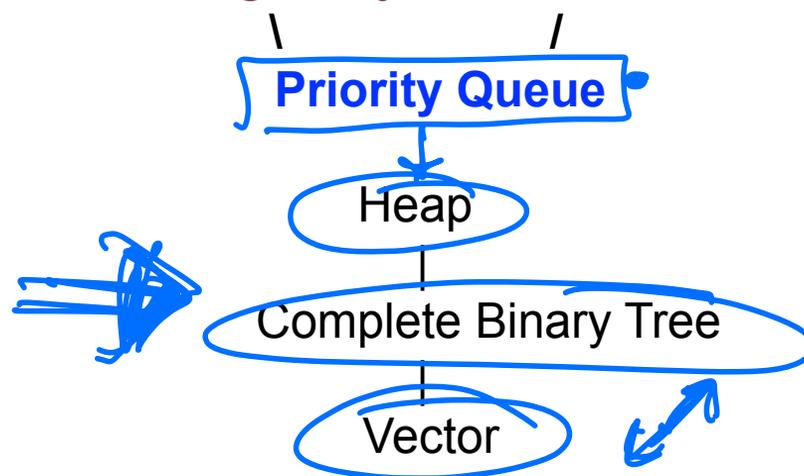
Algorithms: BFS    Task Scheduling

ADT:

Queue

Data structure: Linked list or vector

Huffman Coding    Dijkstra's Shortest Path

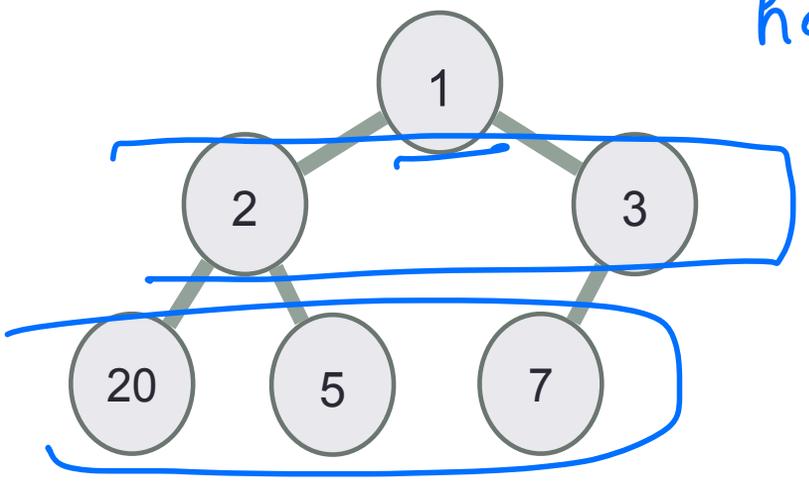


The `priority_queue` **abstract data type (ADT)** is implemented as a **heap**

**Heap is a complete binary tree** with some properties (next lecture)

**Complete Binary Tree** (today!)

# Complete Binary Tree



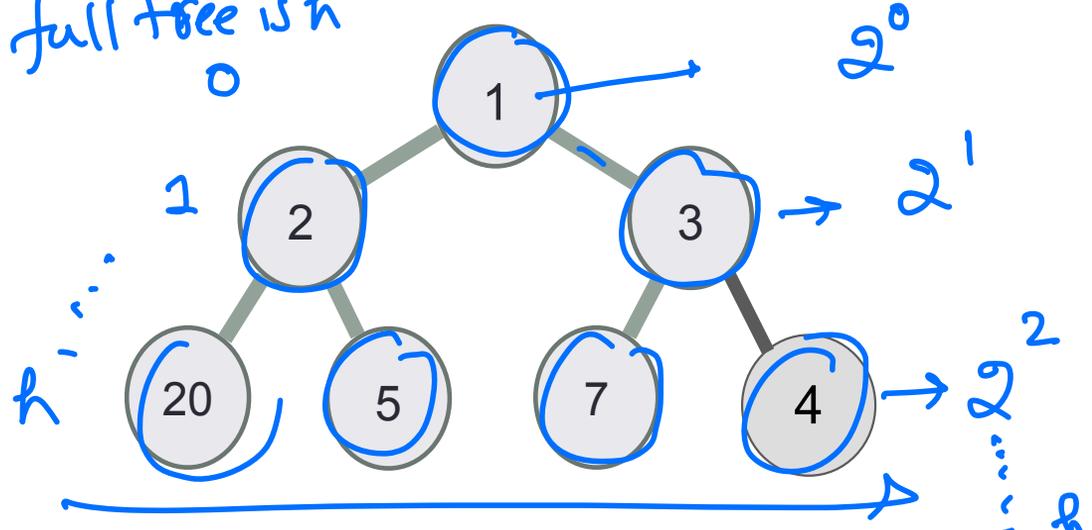
**Complete Binary Tree:**  
 Every level is completely filled (except possibly the last level), and all nodes on the last level are as far left as possible

Complete/full binary trees are **balanced trees!**

$$2^0 + 2^1 + 2^2 + \dots + 2^7 = 2^8 - 1$$

height of full tree is  $h$

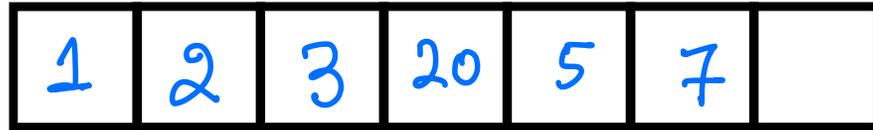
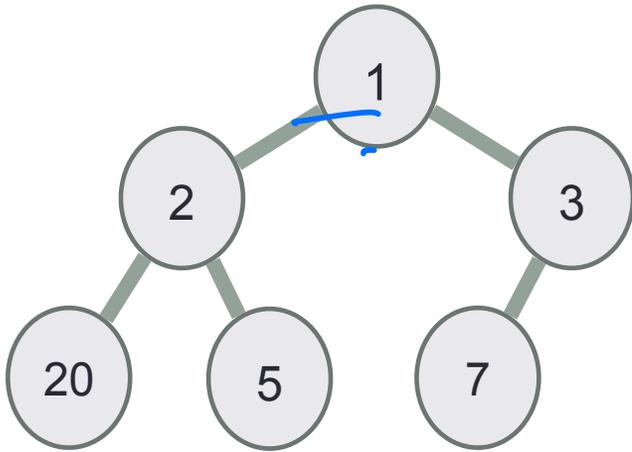
$h$



**Full Binary Tree:** A complete binary tree with last level completely filled

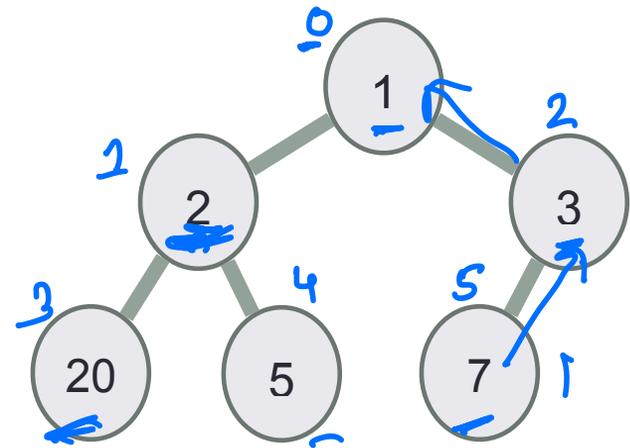
$$2^0 + 2^1 + \dots + 2^h = 2^{h+1} - 1 = n$$

# Representing a complete binary tree as a vector!



- How is the index of each key related to the index of its parent?
- How is the index of each key related to the indices of its left and right child?

# Representing a complete binary tree as a vector!



index
key
parent
left child
right child

1	2	3	20	5	7	
---	---	---	----	---	---	--

0	1	2	3	4	5	
-1	0	0	1	1	2	
1	3	5	/	/	/	
2	4	/	/	/	/	

← indices.

Root is at index 0

For a key at index  $i$ , index of its

- parent is  $\lfloor (i - 1) / 2 \rfloor$
- left child is  $2i + 1$
- right child is  $2i + 2$

$$\lfloor \frac{1-1}{2} \rfloor = 0$$

$$2 \cdot 1 + 1 = 3$$

$$2 \cdot 1 + 2 = 4$$

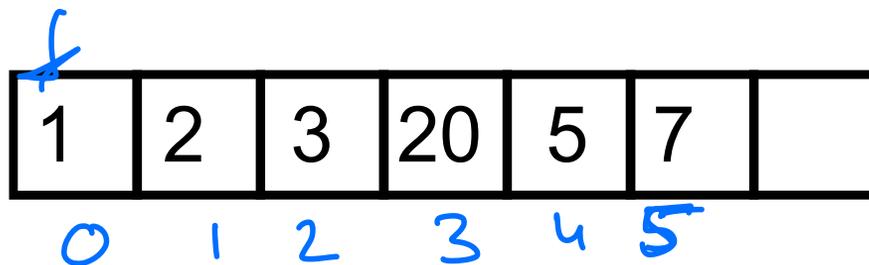
**Activity 2:** For a key at index  $i$ , determine the indices of its parent and children.

# Traverse up the tree using the vector (only)!

Root is at index 0

For a key at index  $i$ , index of its

- parent is  $\lfloor (i - 1) / 2 \rfloor$
- left child is  $2i + 1$
- right child is  $2i + 2$



**Activity 3:** Starting at the last node in the last level (7), write the indices of the keys visited on the path to the root node with key (1):

A. 5, 4, 3, 2, 1, 0

B. 5, 4, 2, 1, 0

C. 5, 3, 1

**D.** 5, 2, 0

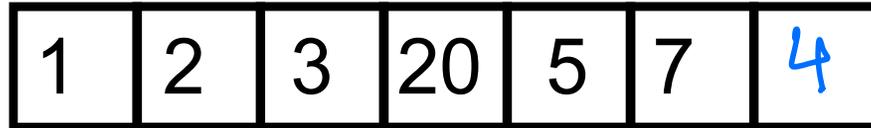
E. None of the above

$$5, \quad \frac{5-1}{2} = 2$$

$$\begin{matrix} 5 \\ (7) \end{matrix}, \quad \begin{matrix} 2 \\ (3) \end{matrix}, \quad \begin{matrix} 0 \\ (1) \end{matrix}$$

Activity 4: Add a new key with value 4 to the complete binary tree represented by the following vector

$O(1)$

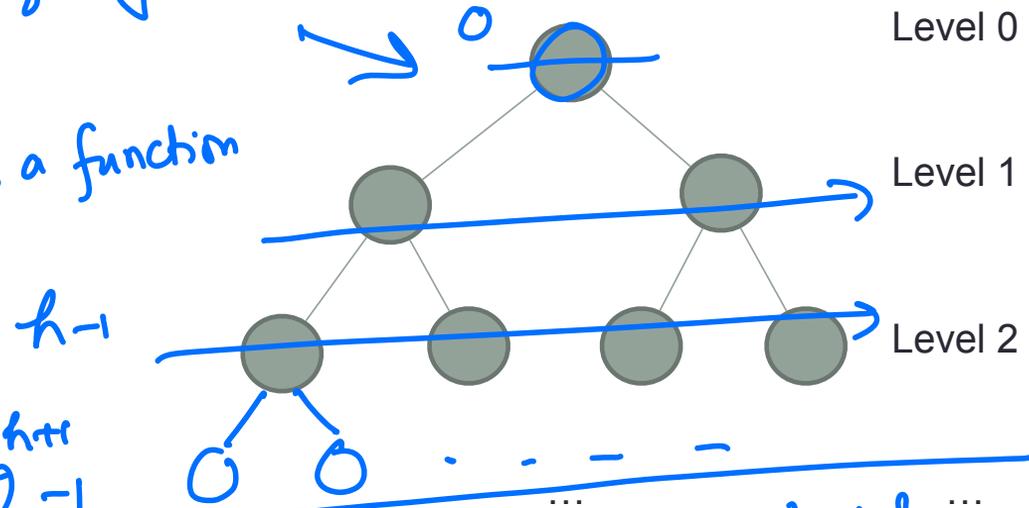


What is the complexity of adding new keys to a complete binary tree?

- A.  $O(1)$
- B.  $O(\log n)$
- C.  $O(n)$
- D. None of the above

# Activity 5: Show that a complete binary tree is balanced

height =  $h$       no. of keys =  $n$   
 To Show       $h = O(\log n)$   
 Approach: Upper bound  $h$  as a function  
 of  $n$



Total no. of keys  
 in levels 0 to  $h-1 \leq n \leq 2^h - 1$

$$2^h - 1 \leq n$$

$$\Rightarrow 2^h \leq n+1$$

$$\Rightarrow h \leq \log(n+1) \Rightarrow h = O(\log n)$$

Total no. of nodes in levels  
 0 to  $h-1$  is:  $2^0 + 2^1 + 2^2 + \dots + 2^{h-1} = 2^h - 1$

## Related Leetcode problems to attempt in problem set 3:

- Level Order Traversal of Binary Tree (medium): <<https://leetcode.com/problems/binary-tree-level-order-traversal/description/?envType=problem-list-v2&envId=binary-tree>>
- Binary Level Order Traversal II (medium): <<https://leetcode.com/problems/binary-tree-level-order-traversal-ii/description/?envType=problem-list-v2&envId=binary-tree>>